

Using Concurrent Function Evaluations to Identify Local Minima of a Derivative-free Optimization Problem

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Motivation

- ▶ We want to identify distinct, “high-quality”, local minimizers of

$$\text{minimize } f(x)$$

$$l \leq x \leq u$$

$$x \in \mathbb{R}^n$$

- ▶ High-quality can be measured by more than the objective.



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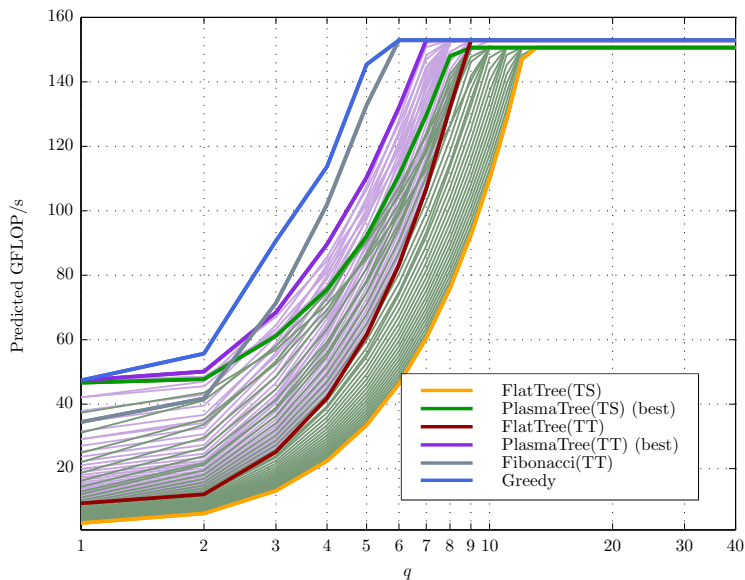
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- ▶ High-quality can be measured by more than the objective.
- ▶ Derivatives of f may or may not be available.
- ▶ The simulation f is likely using parallel resources, but it does not utilize the entire machine.



Why concurrency? Tiled QR example



[Bouwmeester, et al., Tiled QR Factorization Algorithms, 2011]



Global optimization is difficult

Theorem (Törn and Žilinskas, *Global Optimization*, 1989)

An algorithm converges to the global minimum of any continuous f on a domain \mathcal{D} if and only if the algorithm generates iterates that are dense in \mathcal{D} .



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The theory can be more than merely checking that a method generates iterates which are dense in the domain.



Multistart Methods

Given some local optimization routine \mathcal{L} :

Algorithm 1: General Multistart

for $k = 1, 2, \dots$ **do**

 Evaluate f at N points drawn from \mathcal{D}

 Start \mathcal{L} at some set (possibly empty) of previously evaluated points



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- ▶ Which points should start runs?
- ▶ If resources are limited, how should points from each run receive priority?
- ▶ Ideally, only one run is started for each minima.
- ▶ Exploring by sampling. Refining with \mathcal{L} .



Multi-Level Single Linkage

Given some local optimization routine \mathcal{L} :

Algorithm 2: MLSL

for $k = 1, 2, \dots$ **do**

 Sample f at N random points drawn uniformly from \mathcal{D}

 Start \mathcal{L} at all sample points x :

- ▶ that has yet to start a run
 - ▶ $\nexists x_i : \|x - x_i\| \leq r_k$ and $f(x_i) < f(x)$
-

[Rinnooy Kan and Timmer, *Mathematical Programming*, 39(1):57–78, 1987]



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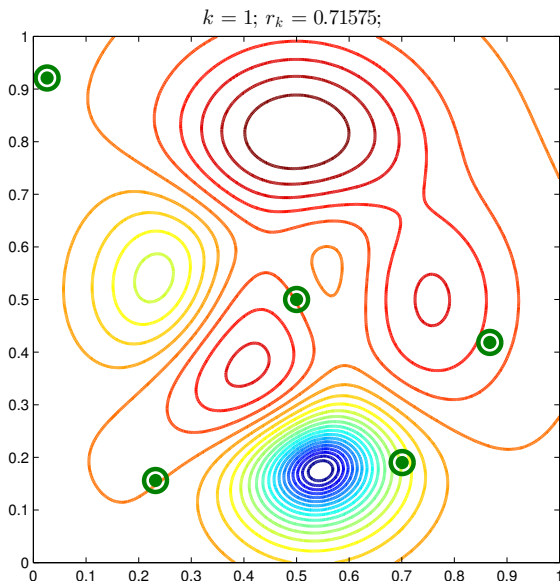
```
for  $k = 1, 2, \dots$  do
    Sample  $f$  at  $N$  random points drawn uniformly from  $\mathcal{D}$ 
    Start  $\mathcal{L}$  at all sample points  $x$ :
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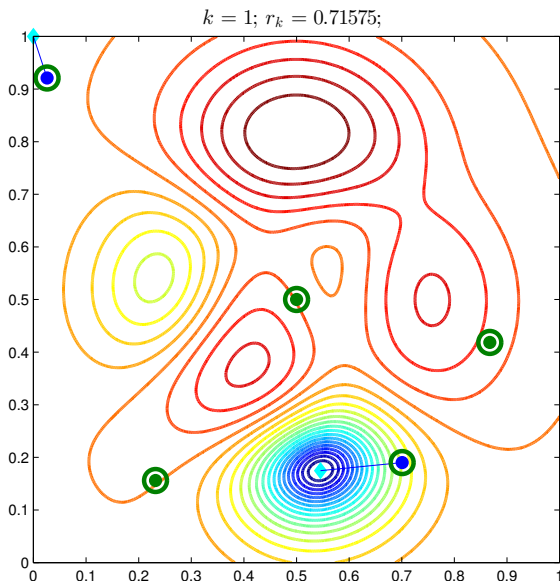
- ▶ Doesn't naturally translate when evaluations of f are limited
- ▶ Ignores some points when deciding where to start \mathcal{L}



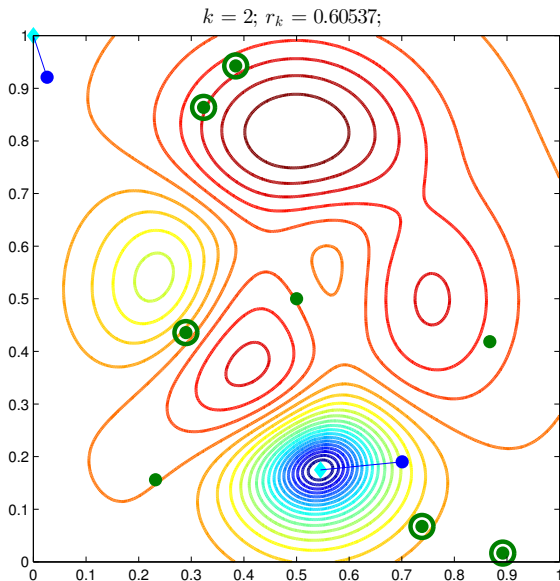
Multi-Level Single Linkage



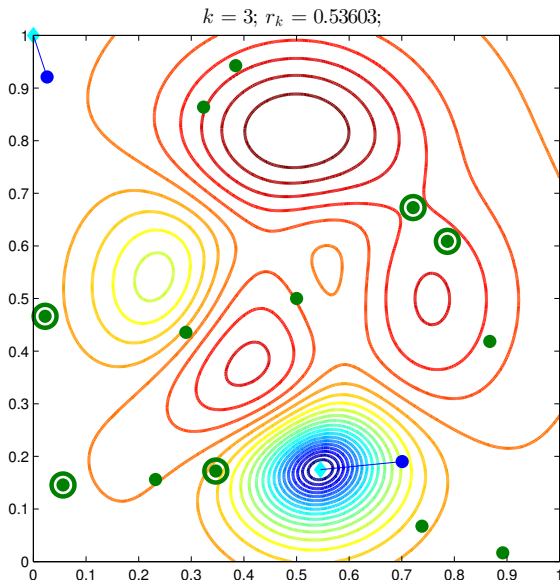
Multi-Level Single Linkage



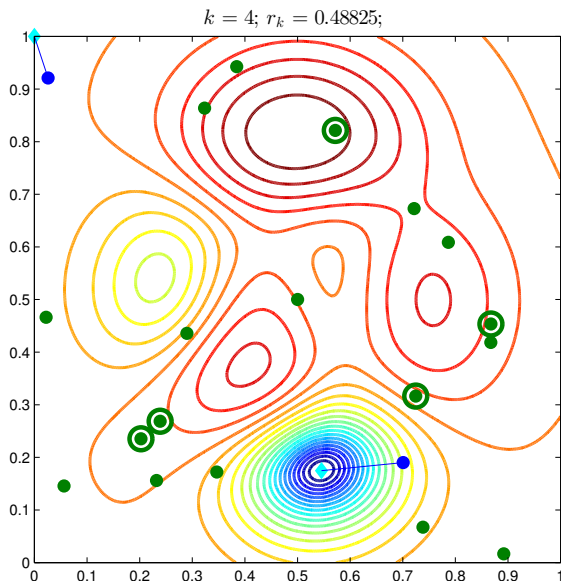
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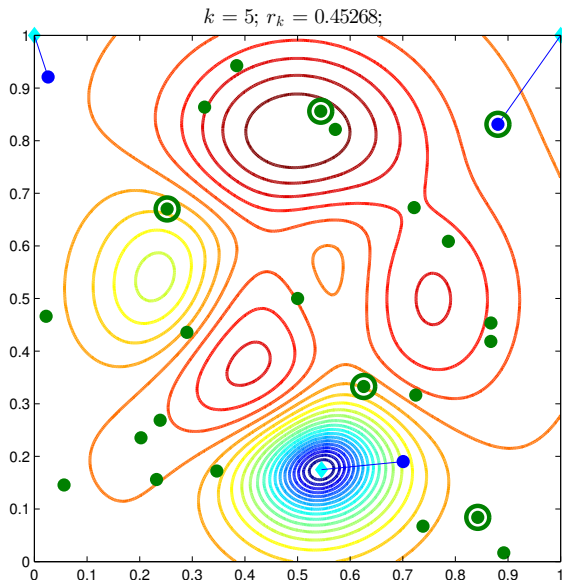
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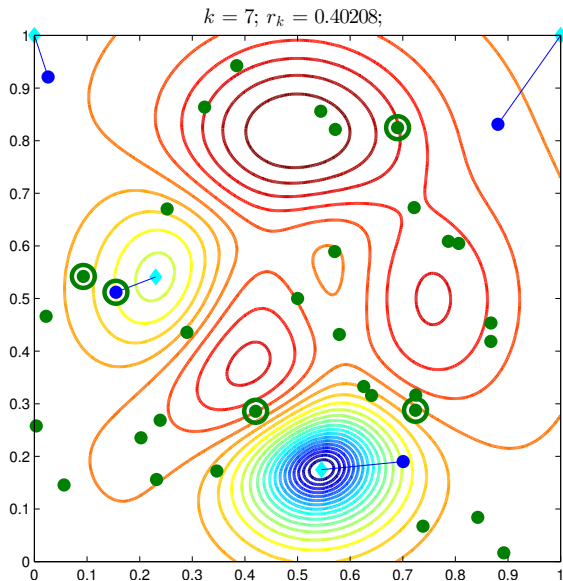
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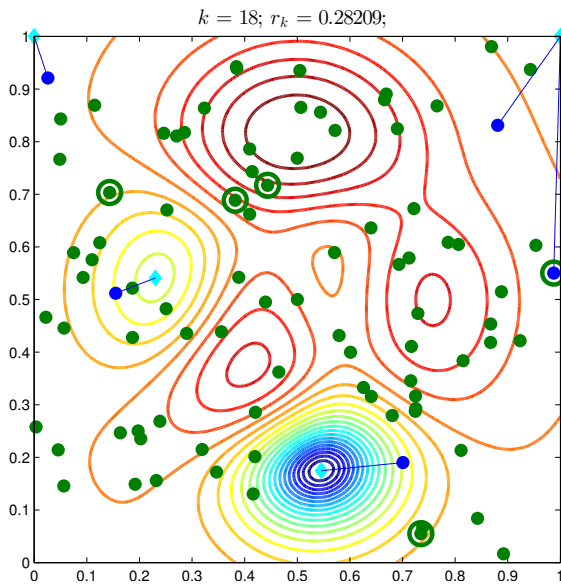
Multi-Level Single Linkage



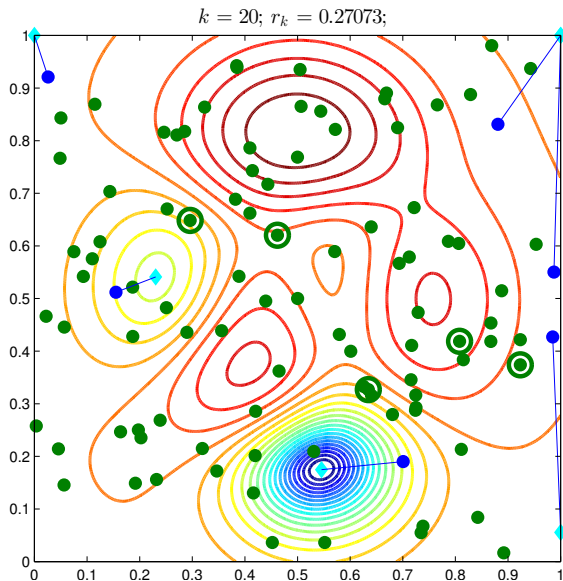
Multi-Level Single Linkage



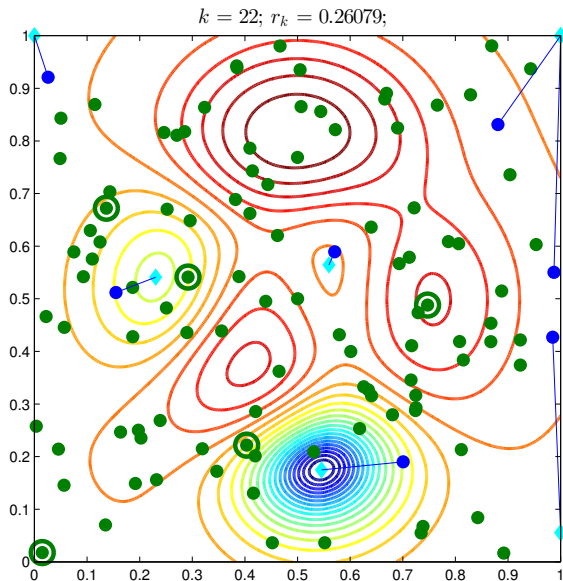
Multi-Level Single Linkage



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Multi-Level Single Linkage

- ▶ $f \in C^2$, with local minima in the interior of \mathcal{D} , and the distance between these minima is bounded away from zero.
- ▶ \mathcal{L} is strictly descent and converges to a minimum (not a stationary point).

▶

$$r_k = \frac{1}{\sqrt{\pi}} \sqrt[n]{\Gamma\left(1 + \frac{n}{2}\right) \text{vol}(\mathcal{D}) \frac{\sigma \log kN}{kN}} \quad (1)$$

Theorem

If $r_k \rightarrow 0$, all local minima will be found almost surely.



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Theorem

If $r_k \rightarrow 0$, all local minima will be found almost surely.

If r_k is defined by (1) with $\sigma > 4$, even if the sampling continues forever, the total number of local searches started is finite almost surely.



BAMLM

MLSL: (S2)–(S4)

$$\hat{x} \in \mathcal{S}_k$$

- (S2) $\nexists x \in \mathcal{S}_k$ *with*
[$\|\hat{x} - x\| \leq r_k$ *and* $f(x) < f(\hat{x})$]
- (S3) \hat{x} *has not started a local optimization run*
- (S4) \hat{x} *is at least μ from $\partial\mathcal{D}$ and ν from known local minima*



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- (S1) $\nexists x \in \mathcal{L}_k$ with
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- (S2) $\nexists x \in \mathcal{S}_k$ with
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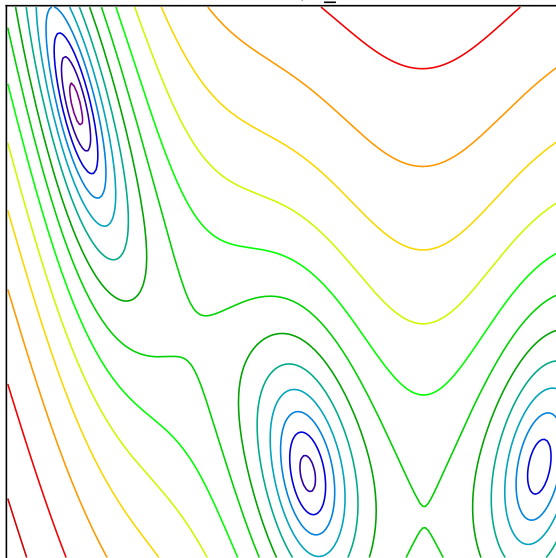
BAMLM: (S1)–(S4), (L1)–(L6)

$$\hat{x} \in \mathcal{L}_k$$

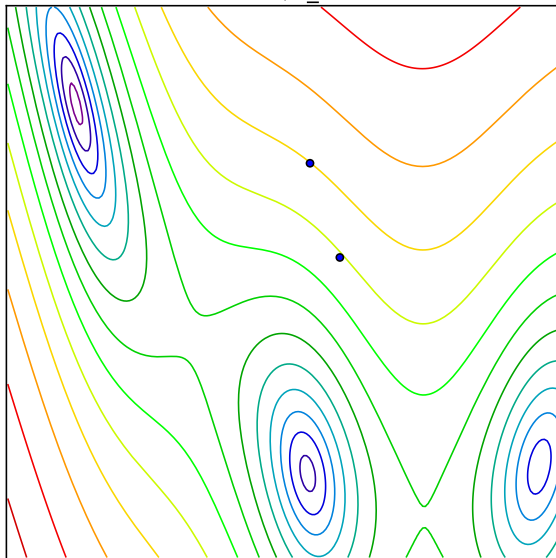
- (L1) $\nexists x \in \mathcal{L}_k$
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- (L4) \hat{x} is at least μ from $\partial\mathcal{D}$ and ν
from known local minima
- (L5) \hat{x} is not in an active local
optimization run and has not
been ruled stationary
- (L6) $\exists r_k$ -descent path in \mathcal{H}_k from
some $x \in \mathcal{S}_k$ satisfying (S2–S4)
to \hat{x}



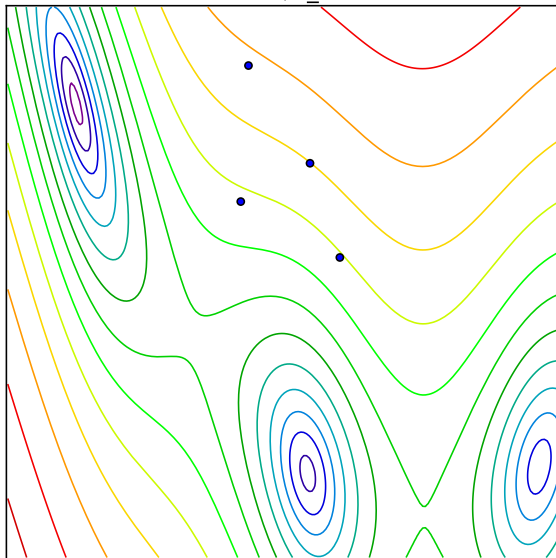
Iteration: 0; r_k : Inf



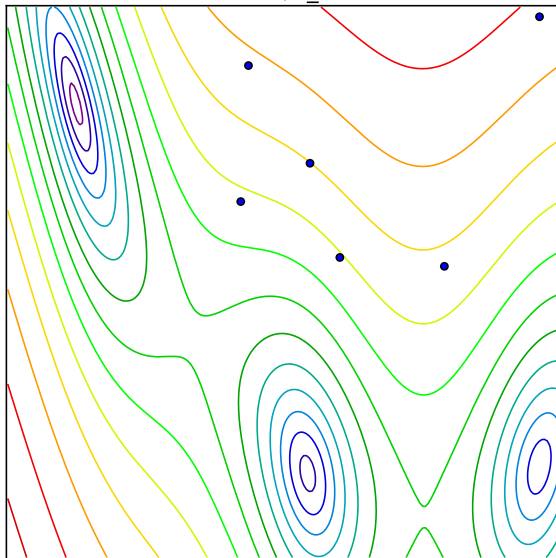
Iteration: 1; r_k : 0.743



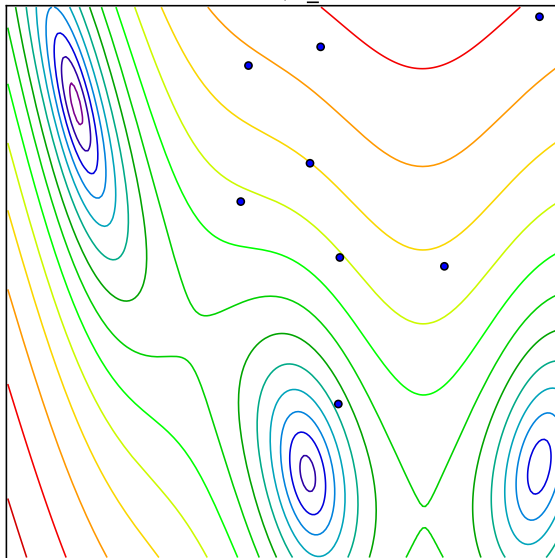
Iteration: 2; r_k : 0.743



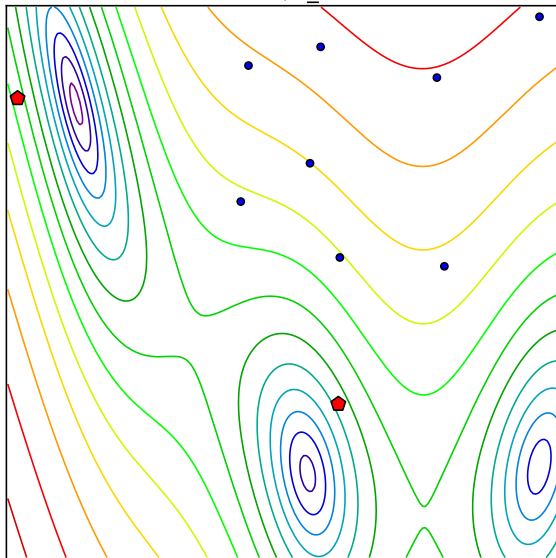
Iteration: 3; r_k : 0.689



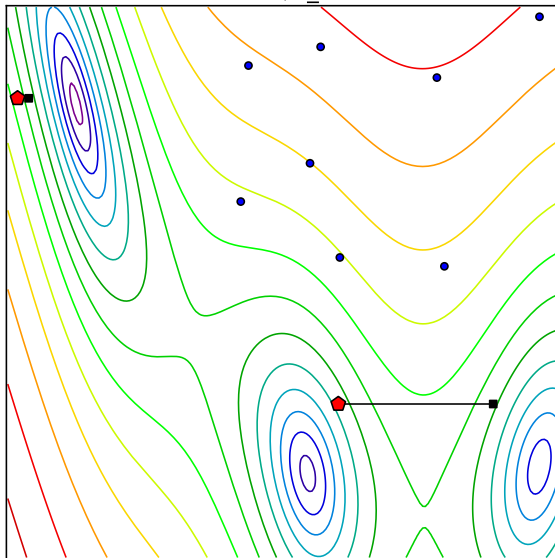
Iteration: 4; r_k : 0.643



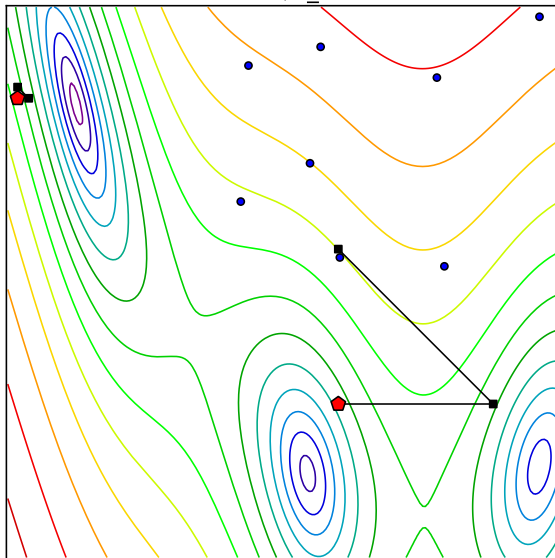
Iteration: 5; r_k : 0.605



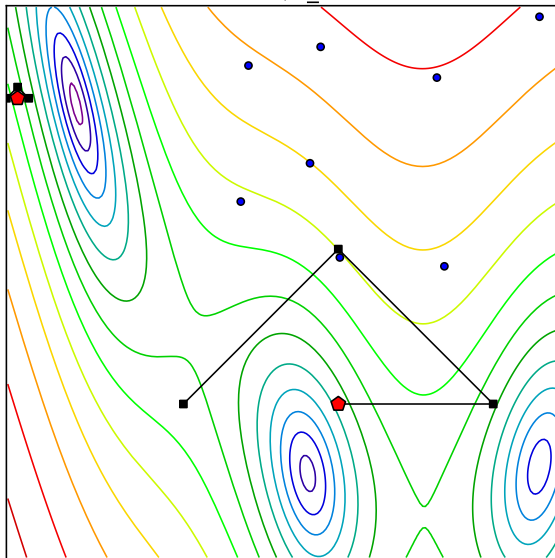
Iteration: 6; r_k : 0.605



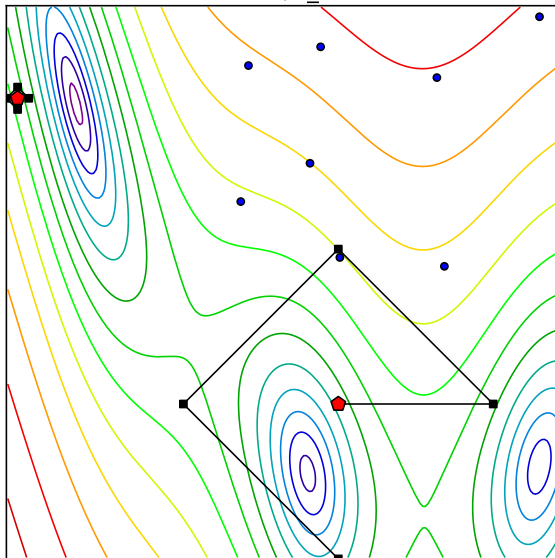
Iteration: 7; r_k : 0.605



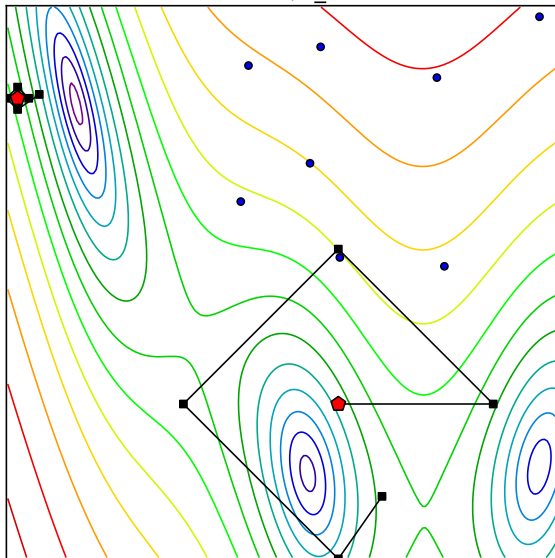
Iteration: 8; r_k : 0.605



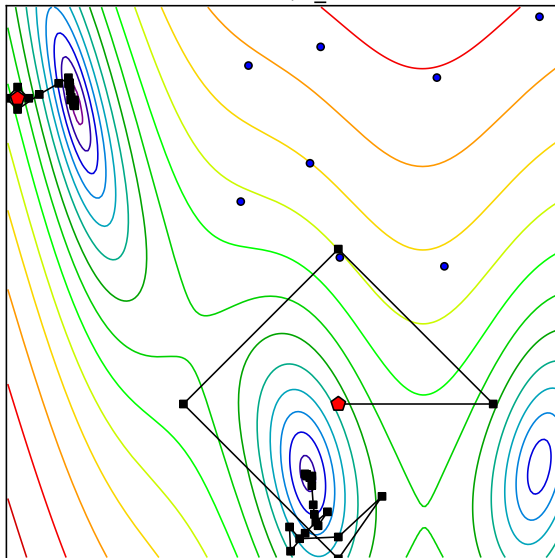
Iteration: 9; r_k : 0.605



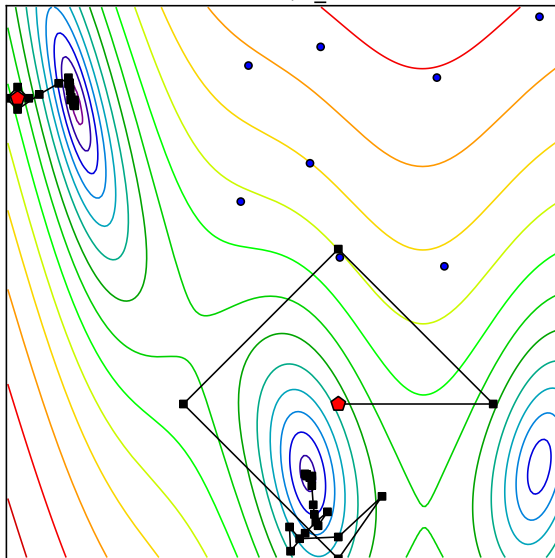
Iteration: 10; r_k : 0.605



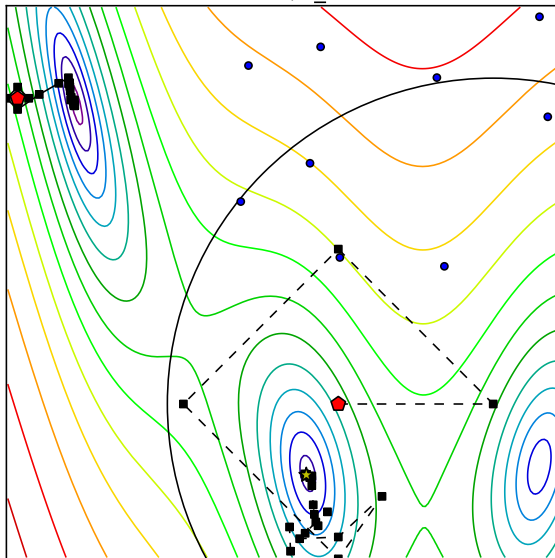
Iteration: 35; r_k : 0.605



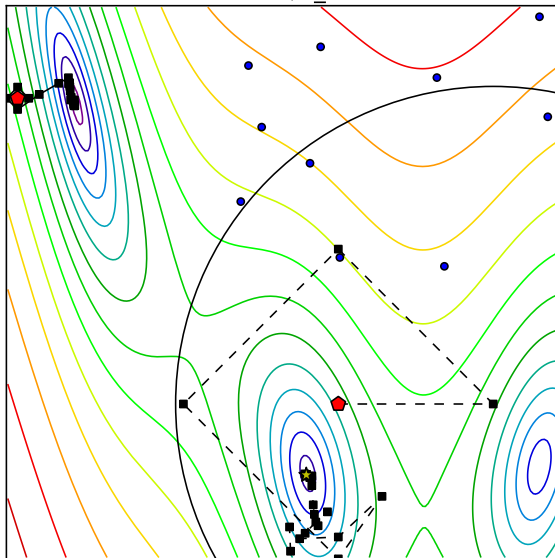
Iteration: 36; r_k : 0.605



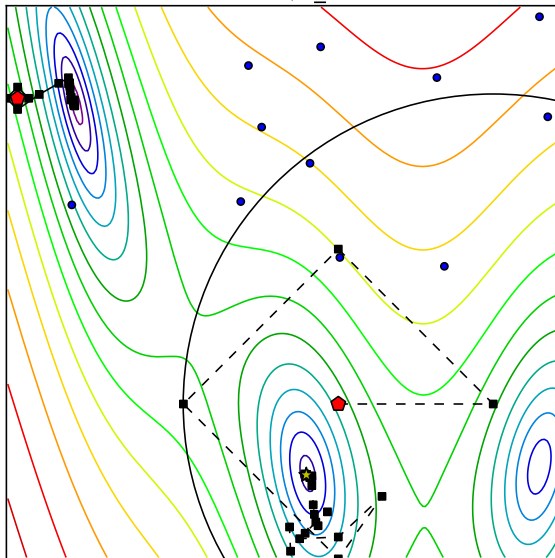
Iteration: 37; r_k : 0.589



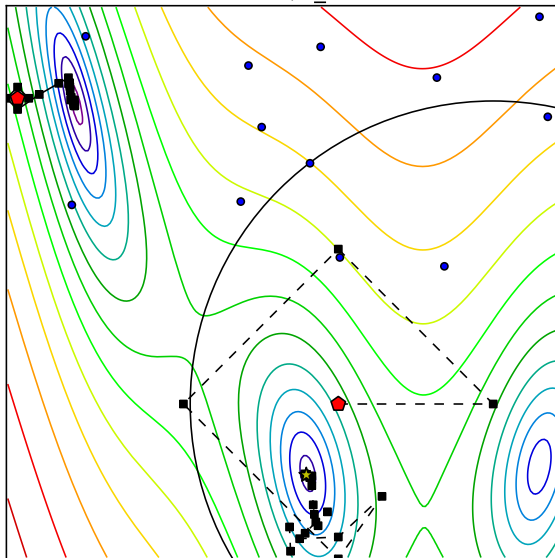
Iteration: 38; r_k : 0.574



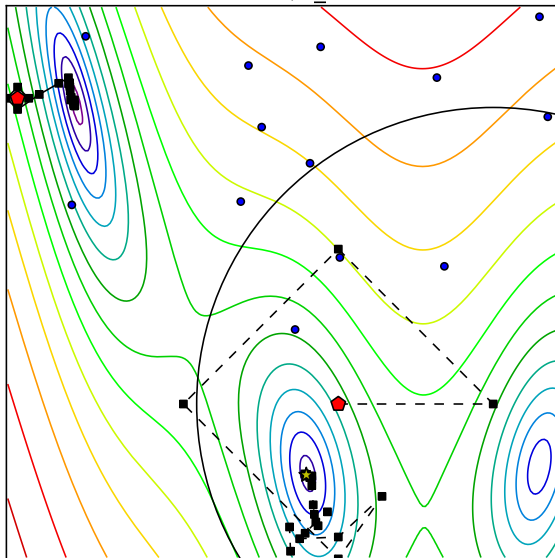
Iteration: 39; r_k : 0.560



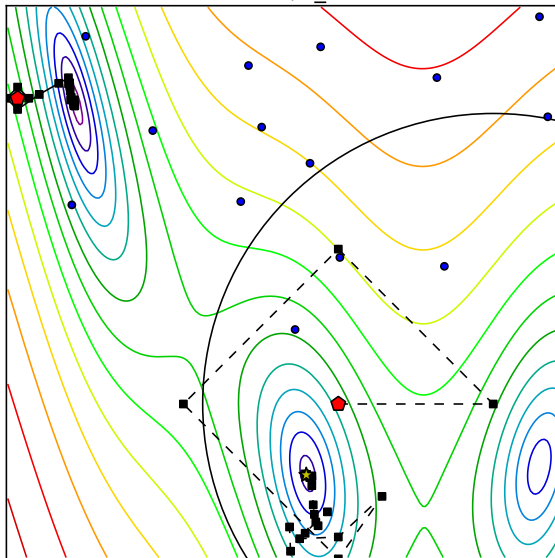
Iteration: 40; r_k : 0.548



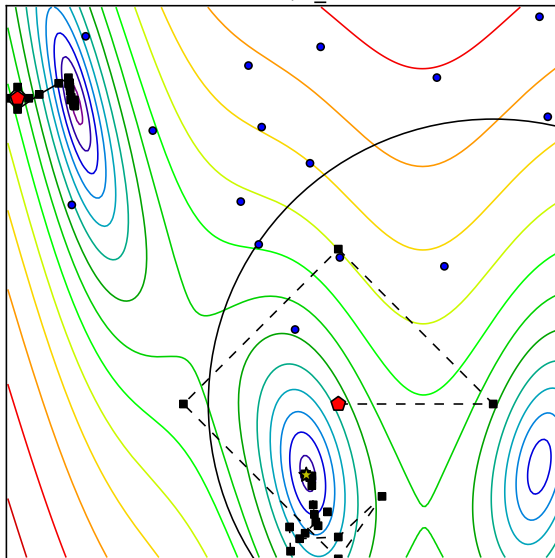
Iteration: 41; r_k : 0.536



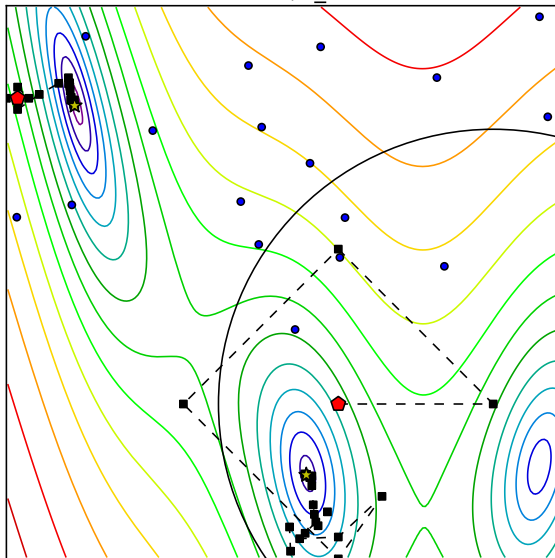
Iteration: 42; r_k : 0.525



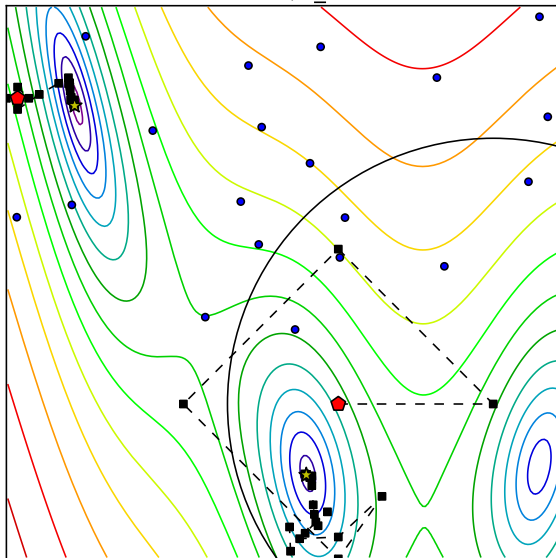
Iteration: 43; r_k : 0.515



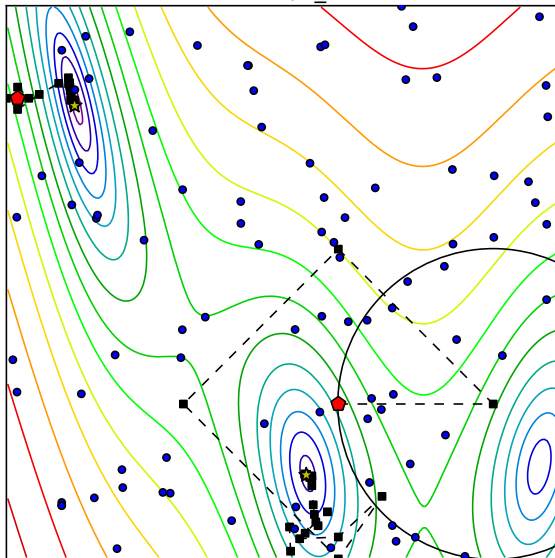
Iteration: 44; r_k : 0.497



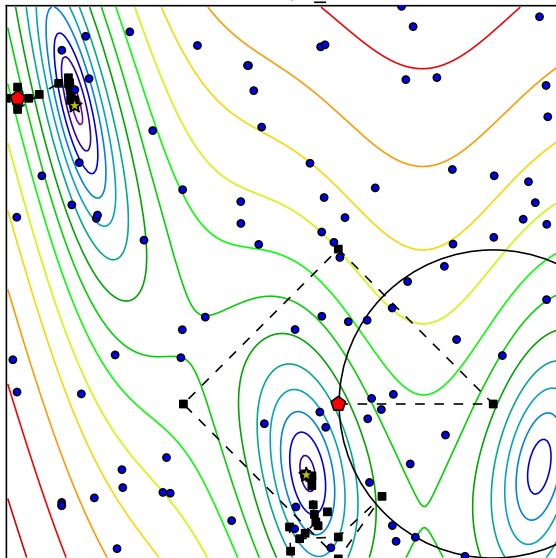
Iteration: 45; r_k : 0.480



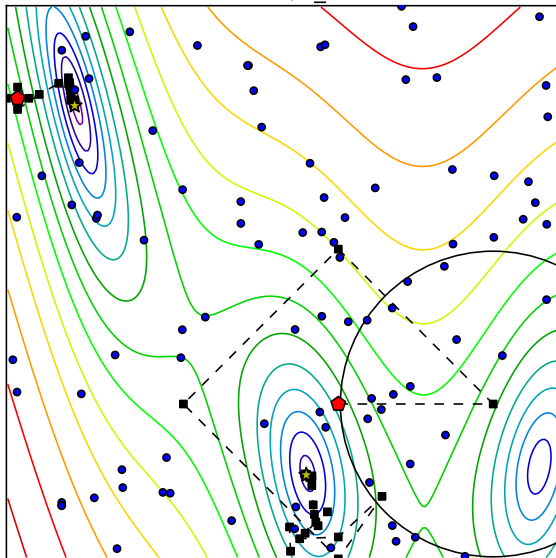
Iteration: 80; r_k : 0.281



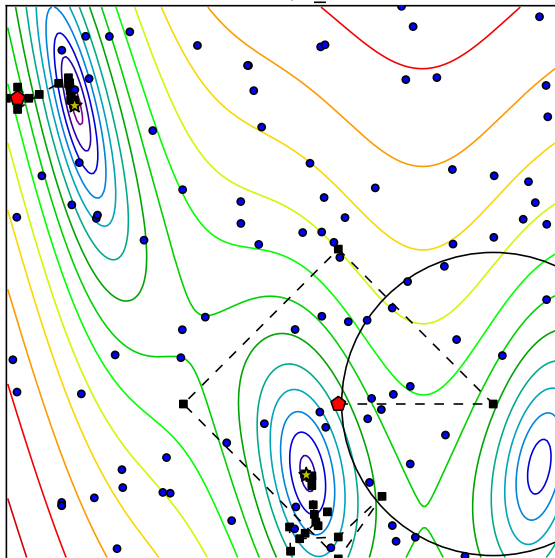
Iteration: 81; r_k : 0.279



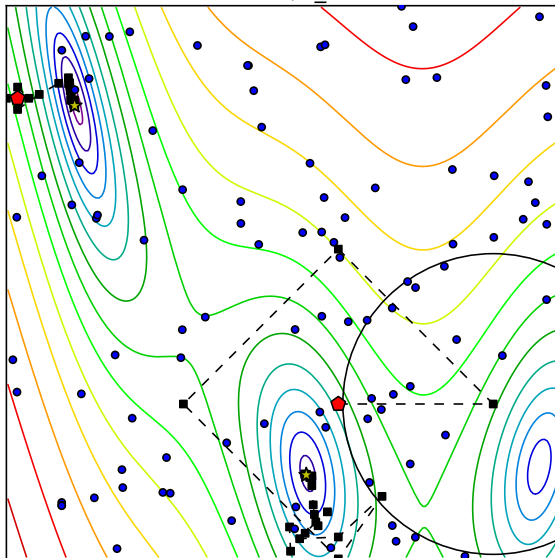
Iteration: 82; r_k : 0.276



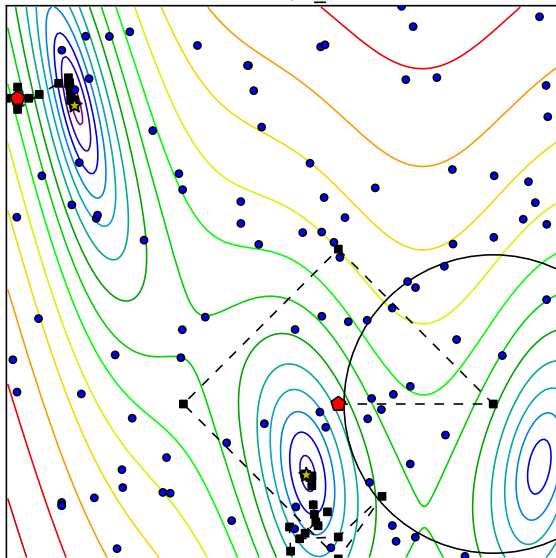
Iteration: 83; r_k : 0.274



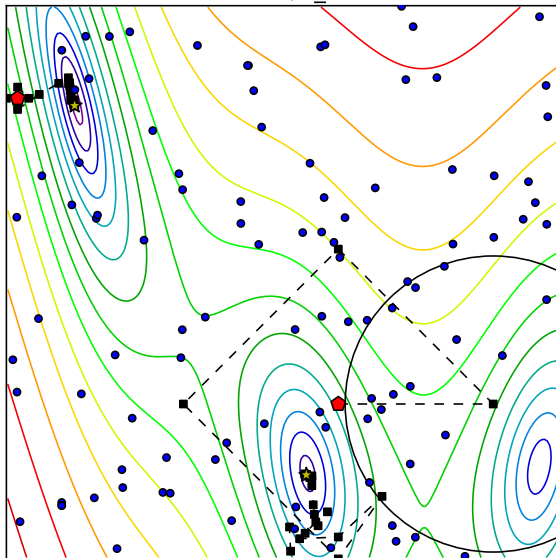
Iteration: 84; r_k : 0.272



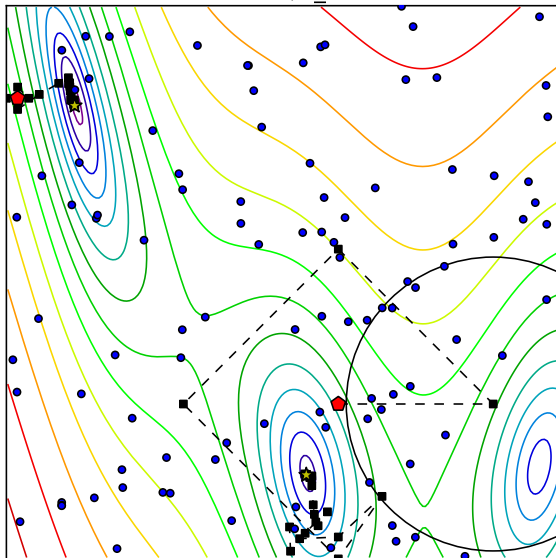
Iteration: 85; r_k : 0.270



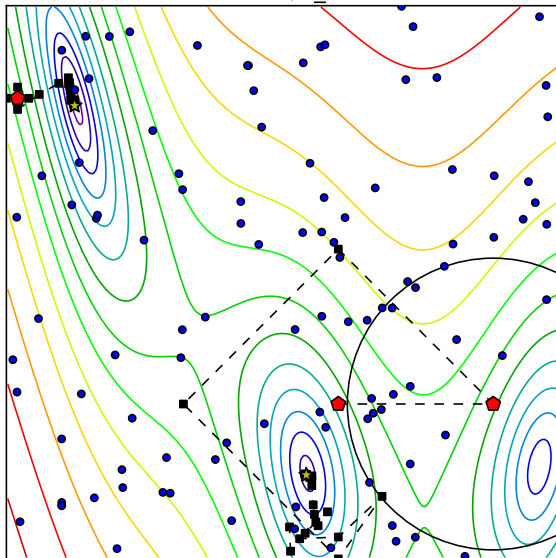
Iteration: 86; r_k : 0.268



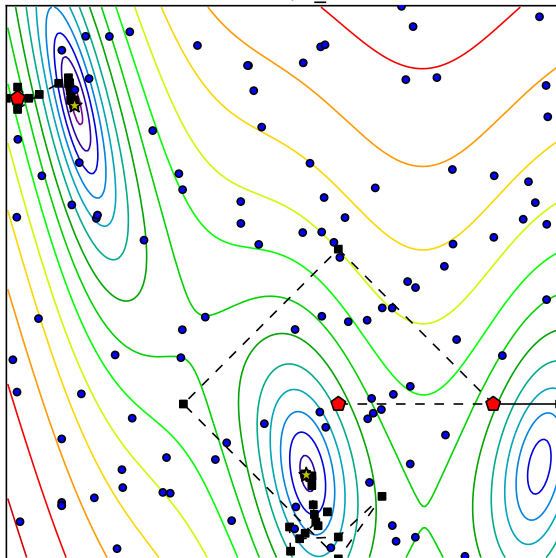
Iteration: 87; r_k : 0.266



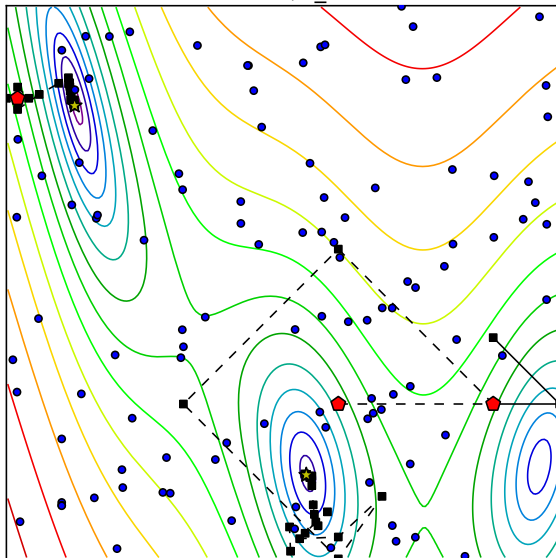
Iteration: 88; r_k : 0.264



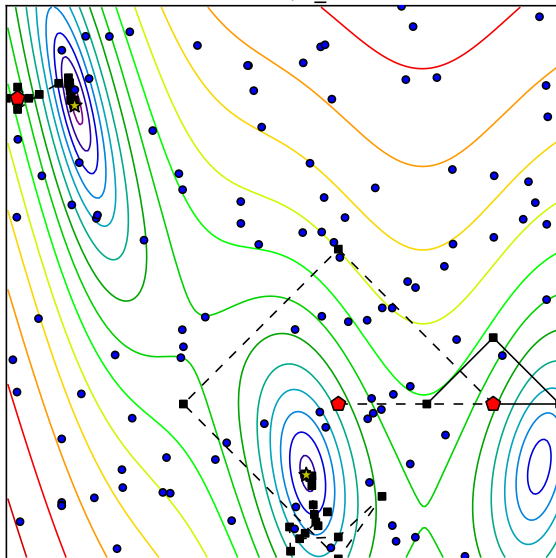
Iteration: 89; r_k : 0.263



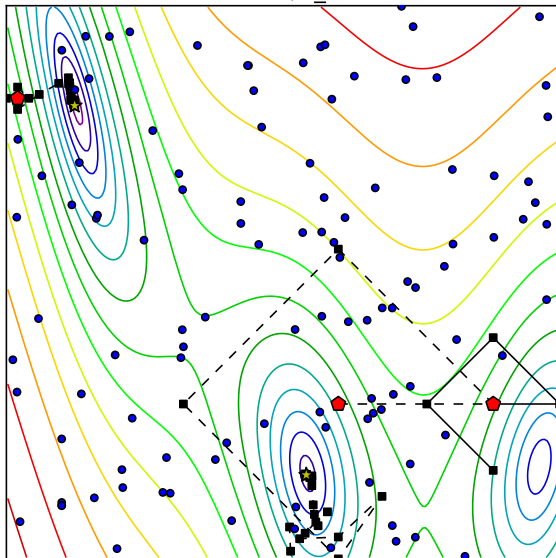
Iteration: 90; r_k : 0.262



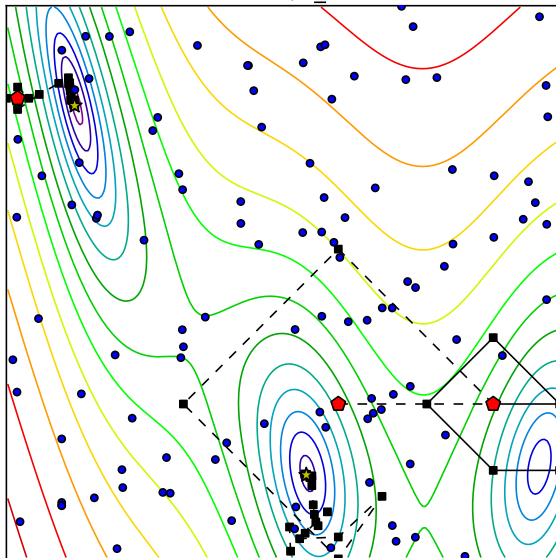
Iteration: 91; r_k : 0.261



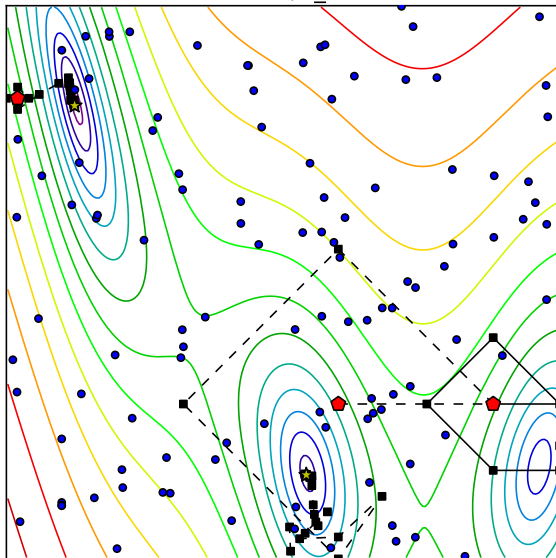
Iteration: 92; r_k : 0.260



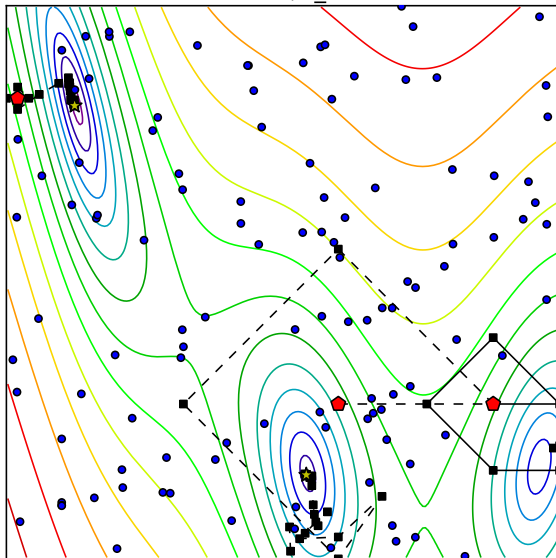
Iteration: 93; r_k : 0.259



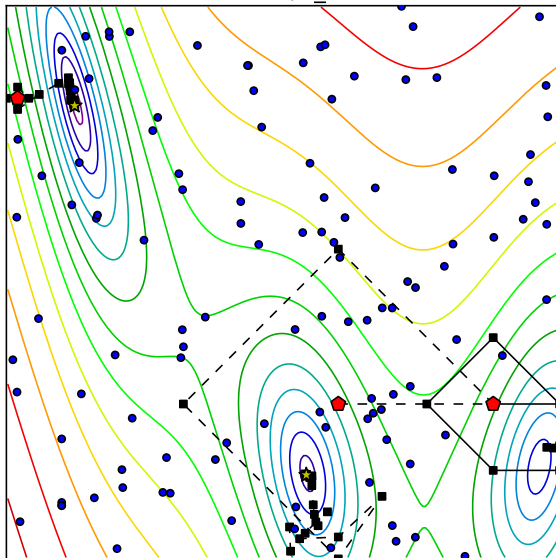
Iteration: 94; r_k : 0.258



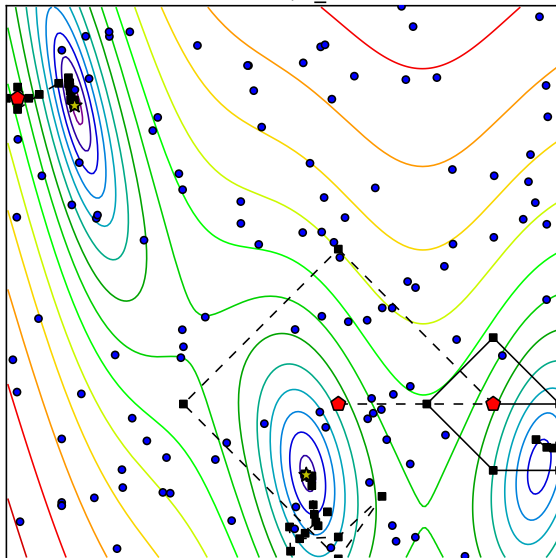
Iteration: 95; r_k : 0.257



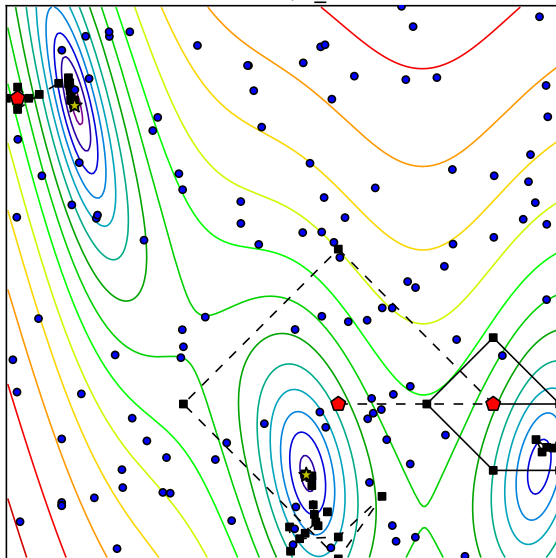
Iteration: 96; r_k : 0.256



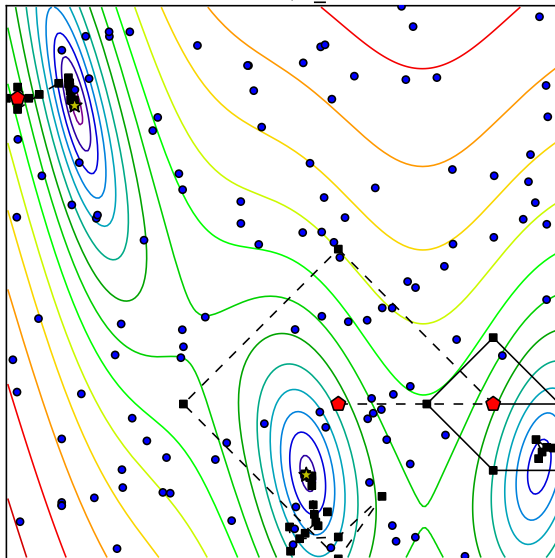
Iteration: 97; r_k : 0.255



Iteration: 98; r_k : 0.255



Iteration: 99; r_k : 0.254



Properties of the local optimization method

Necessary:

- ▶ Honors a starting point
- ▶ Honors bound constraints



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ORBIT satisfies these [Wild, Regis, Shoemaker, SIAM-JOSC, 2008]

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Possibly beneficial:

- ▶ Can return multiple points of interest
- ▶ Reports solution quality/confidence at every iteration
- ▶ Can avoid certain regions in the domain
- ▶ Uses a history of past evaluations of f
- ▶ Uses additional points mid-run



Algorithm 3: AAML

Give each worker a point to evaluate

for $k = 1, 2, \dots$ **do**

 Receive from (longest waiting) worker w that has evaluated f

 Update \mathcal{H}_k and r_k

if *point evaluated by w is from an active run* **then**

if *Run is complete* **then**

 Update X_k^* , and mark points inactive

else

 Add the next point in its localopt run (not in \mathcal{H}_k) to Q_L

 Start run(s) at all point(s) satisfying (S1)–(S4), (L1)–(L6)

 Add the subsequent point (not in \mathcal{H}_k) from each run to Q_L

 Merge runs in Q_L with candidate minima within 2ν of each other

 Give w a point at which to evaluate f , either from Q_L or \mathcal{R}

MLSL: (S2)–(S4)

$$\hat{x} \in \mathcal{S}_k$$

- (S1) $\nexists x \in \mathcal{L}_k$ with
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (S2) $\nexists x \in \mathcal{S}_k$ with
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (S3) \hat{x} has not started a local optimization run
- (S4) \hat{x} is at least μ from $\partial\mathcal{D}$ and ν from known local minima

BAMLM: (S1)–(S4), (L1)–(L6)

$$\hat{x} \in \mathcal{L}_k$$

- (L1) $\nexists x \in \mathcal{L}_k$
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (L2) $\nexists x \in \mathcal{S}_k$ with
[$\|\hat{x} - x\| \leq r_k$ and $f(x) < f(\hat{x})$]
- (L3) \hat{x} has not started a local optimization run
- (L4) \hat{x} is at least μ from $\partial\mathcal{D}$ and ν from known local minima
- (L5) \hat{x} is not in an active local optimization run and has not been ruled stationary
- (L6) $\exists r_k$ -descent path in \mathcal{H}_k from some $x \in \mathcal{S}_k$ satisfying (S2-S4) to \hat{x}



Theorem

Given the same assumptions as MLSL, AAML will start a finite number of local optimization runs with probability 1.



AAML Theory

Theorem

Given the same assumptions as MLSL, AAML will start a finite number of local optimization runs with probability 1.

Assumption

There exists $K_0 < \infty$ so that for any K_0 consecutive iterations, there is a positive (bounded away from zero) probability of evaluating a point from the sample stream and each existing local optimization run.



AAML Theory

Theorem

Given the same assumptions as MLSL, AAML will start a finite number of local optimization runs with probability 1.

Assumption

There exists $K_0 < \infty$ so that for any K_0 consecutive iterations, there is a positive (bounded away from zero) probability of evaluating a point from the sample stream and each existing local optimization run.

Theorem

Each $x^* \in X^*$ will almost surely be either identified in a finite number of evaluations or have a single local optimization run that is converging asymptotically to it.



Measuring Performance

GLODS Global & local optimization using direct search [Custódio, Madeira (JOGO, 2014)]

Direct Serial DIRECT [D. Finkel's MATLAB code]

pVTDirect Parallel DIRECT [He, Watson, Sosonkina (TOMS, 2009)]

Random Uniform sampling over domain (as a baseline)

BAMLM

- ▶ Concurrency: 4
- ▶ Local optimization method
 - ▶ ORBIT [Wild, Regis, & Shoemaker (SIAM JOSOC, 2008)]
 - ▶ BOBYQA [Powell, 2009]
- ▶ Initial sample size: $10n$

- ▶ Each method evaluates Direct's $2n + 1$ initial points.



Measuring Performance

Notation:

Let X^* be the set of all local minima of f .

Let $f_{(i)}^*$ be the i th smallest value $\{f(x^*) | x^* \in X^*\}$.

Let $x_{(i)}^*$ be the element of X^* corresponding to the value $f_{(i)}^*$.

The global minimum has been found at a level $\tau > 0$ at batch k if an algorithm it has found a point \hat{x} satisfying:

$$f(\hat{x}) - f_{(1)}^* \leq (1 - \tau) \left(f(x_0) - f_{(1)}^* \right),$$

where x_0 is the starting point for problem p .



Measuring Performance

Notation:

Let X^* be the set of all local minima of f .

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Let $x_{(i)}^*$ be the element of X^* corresponding to the value $f_{(i)}^*$.

The j best local minima have been found at a level $\tau > 0$ at batch k if:

$$\left| \left\{ x_{(1)}^*, \dots, x_{(\underline{j}-1)}^* \right\} \cap \left\{ x_{(i)}^* : \exists x \in \mathcal{H}_k \text{ with } \|x - x_{(i)}^*\| \leq r_n(\tau) \right\} \right| = \underline{j} - 1$$

&

$$\left| \left\{ x_{(\underline{j})}^*, \dots, x_{(\bar{j})}^* \right\} \cap \left\{ x_{(i)}^* : \exists x \in \mathcal{H}_k \text{ with } \|x - x_{(i)}^*\| \leq r_n(\tau) \right\} \right| \geq \bar{j} - \underline{j} + 1,$$

where \underline{j} and \bar{j} are the smallest and largest integers such that

$$f_{(\underline{j})}^* = f_{(\underline{j})}^* = f_{(\underline{j})}^* \text{ and where } r_n(\tau) = \sqrt[n]{\frac{\tau \text{vol}(\mathcal{D}) \Gamma(\frac{n}{2} + 1)}{\pi^{n/2}}}.$$



Problems considered

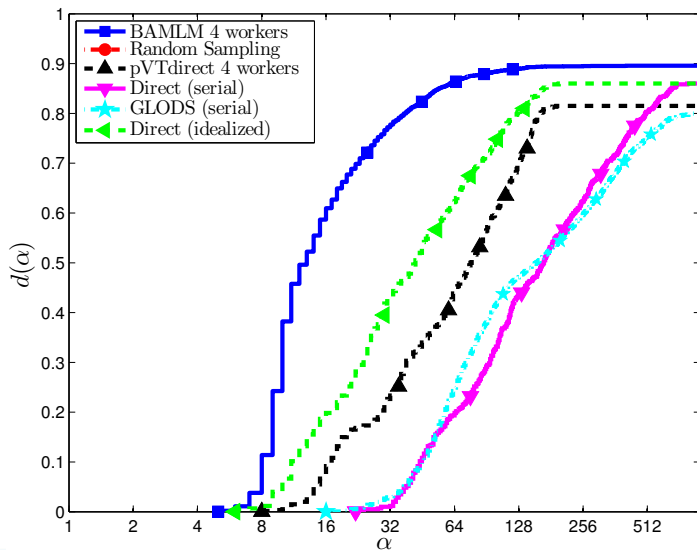
GKLS problem generator [Gaviano et al., “Algorithm 829” (TOMS, 2003)]

- ▶ 600 synthetic problems with known local minima
- ▶ $n = 2, \dots, 7$
- ▶ 10 local minima in the unit cube with a unique global minimum
- ▶ 100 problems for each dimension
- ▶ 5 replications (different seeds) for each problem
- ▶ 5000 evaluations



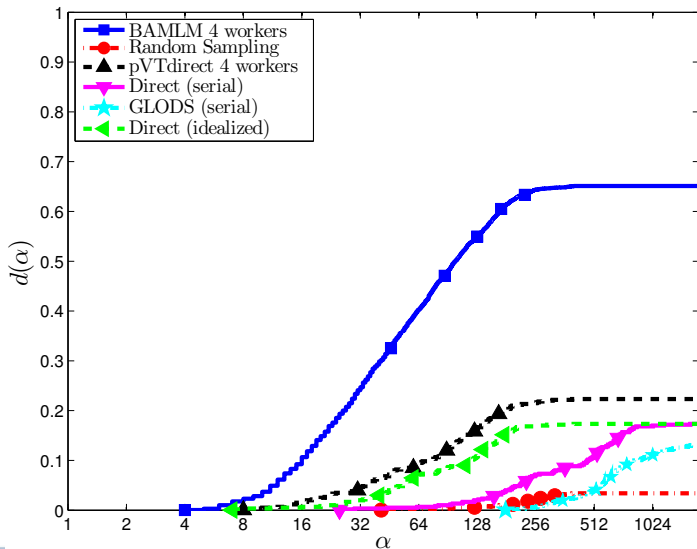
Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



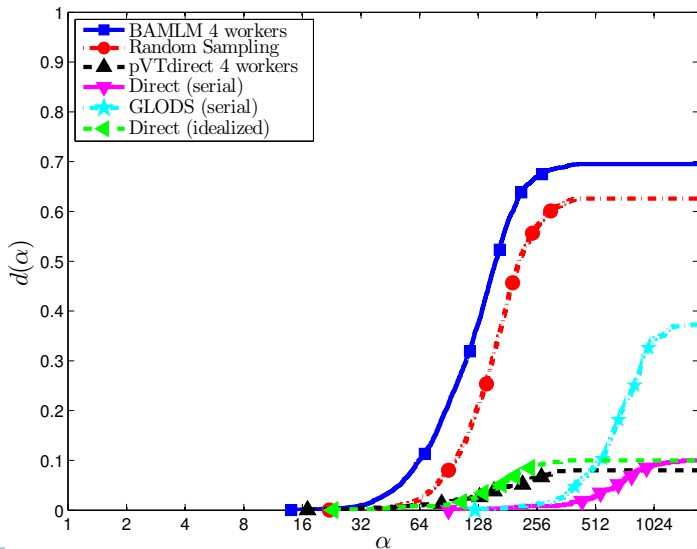
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



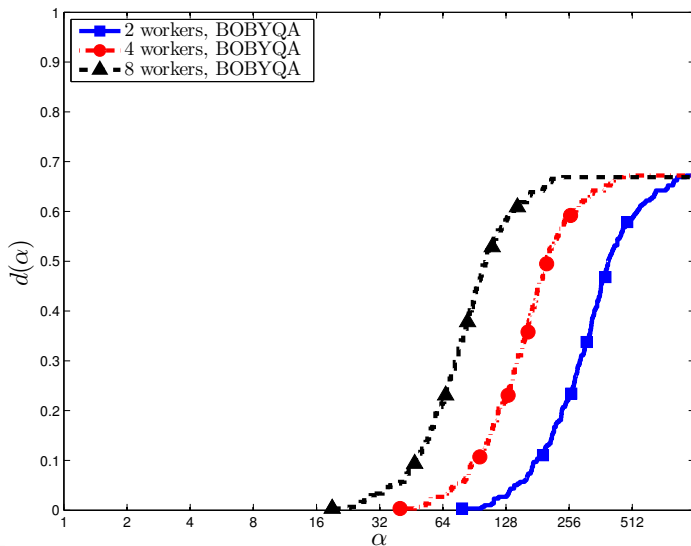
Data Profiles

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 7 best minima



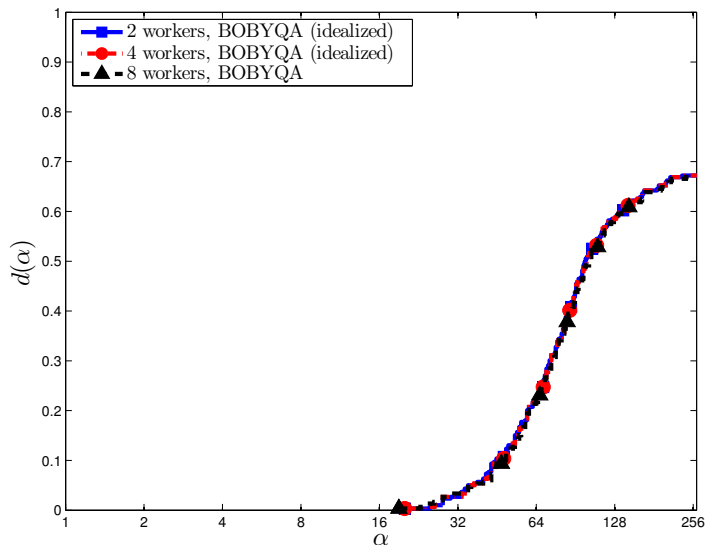
Data Profiles

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 7 best minima



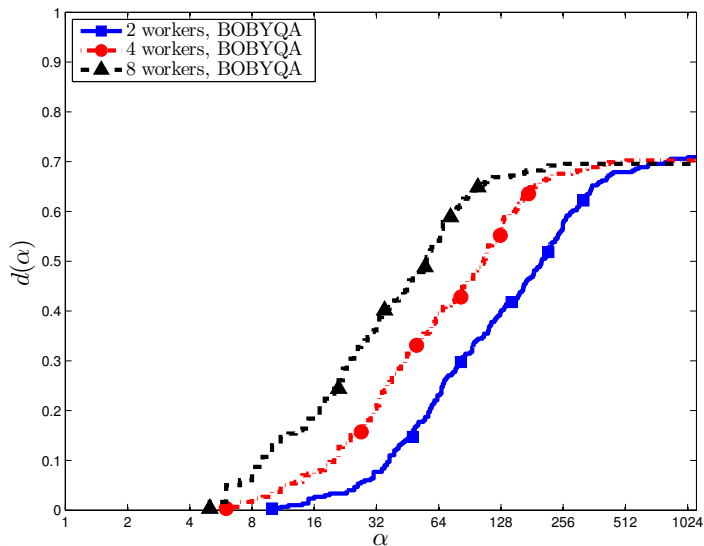
Data Profiles

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 7 best minima



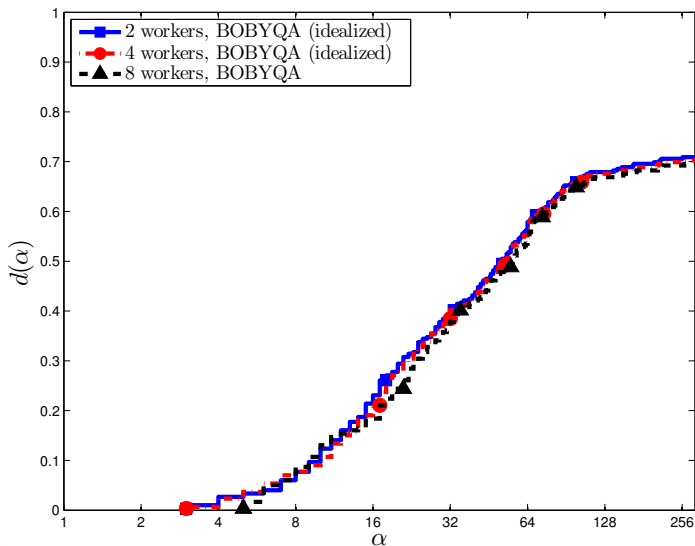
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



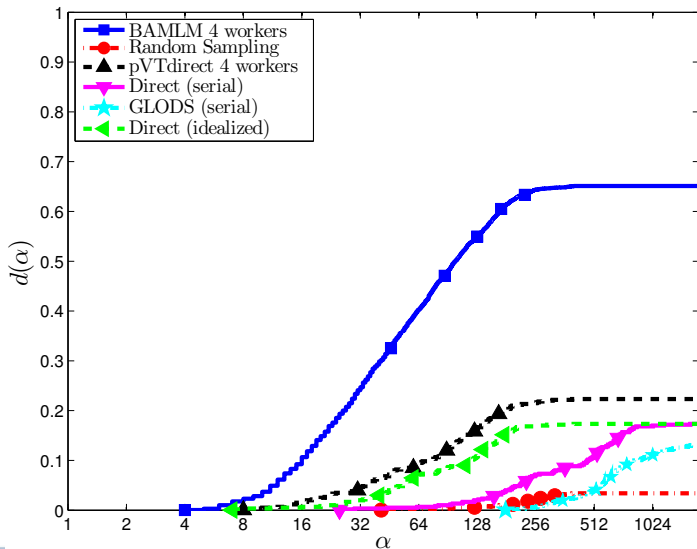
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



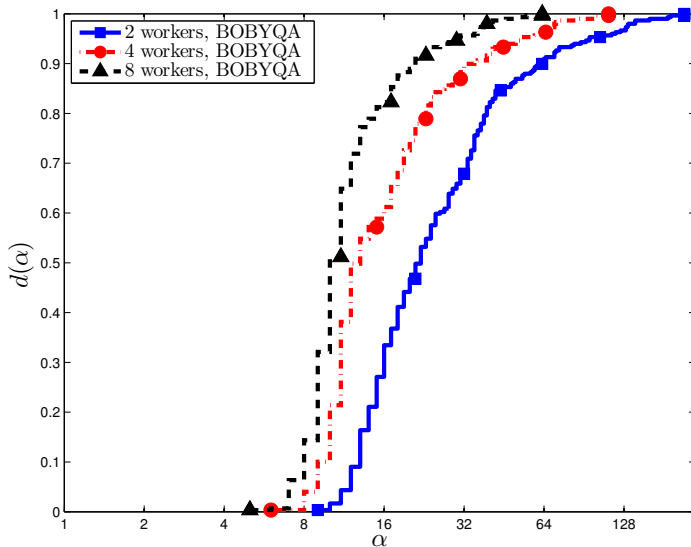
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 3 best minima



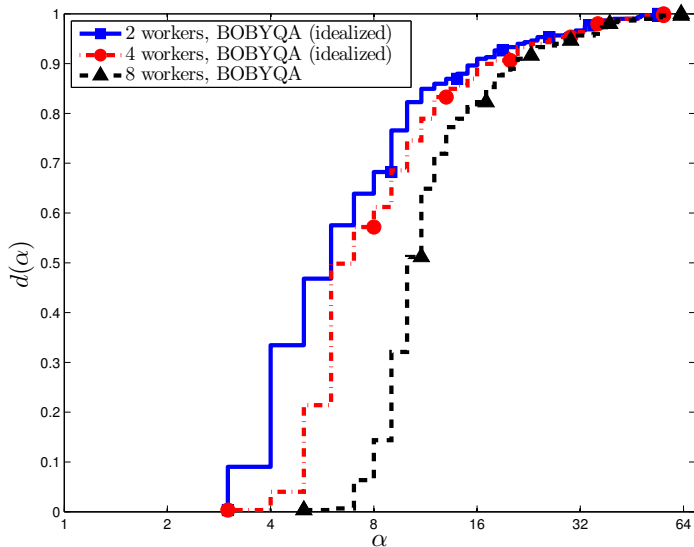
Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$

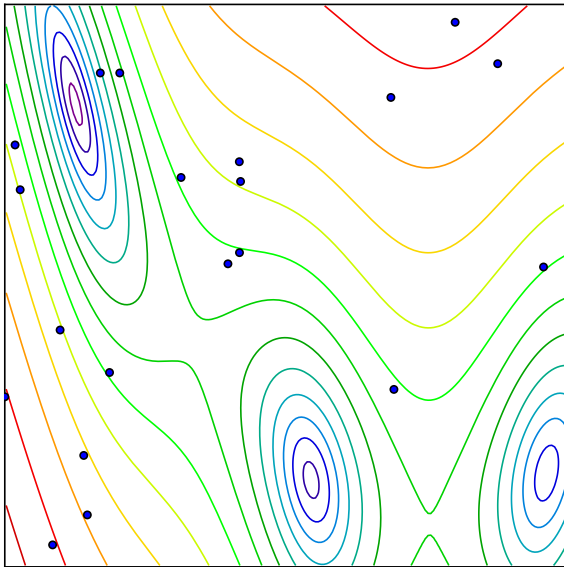


Data Profiles

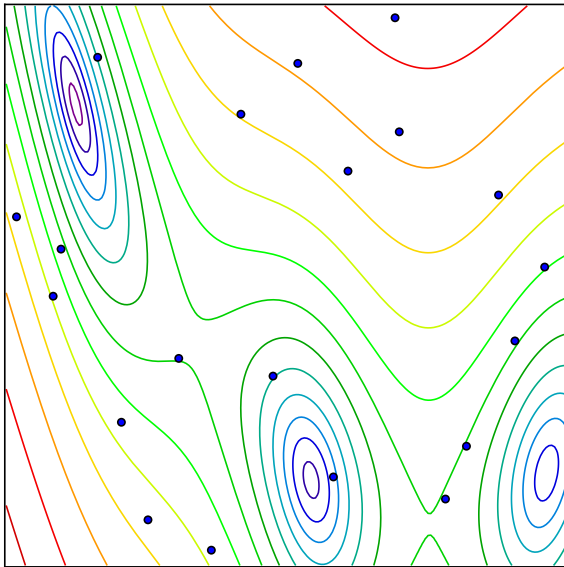
$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



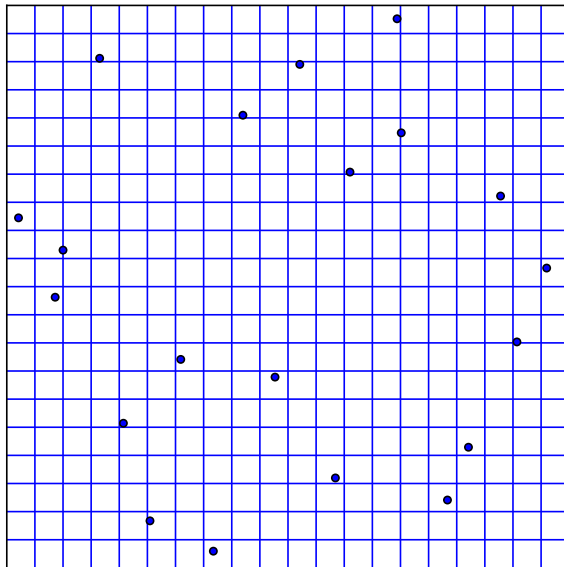
Uniform sampling



Latin hypercube sampling



Latin hypercube sampling



BAMLM with LHS

Critical distance for uniform sampling:

$$r_k = \pi^{-1/2} \left(\Gamma(1 + \frac{n}{2}) \text{vol}(\mathcal{D}) \frac{\sigma \log kN}{kN} \right)^{1/n}$$

Critical distance for Latin hypercube sampling:

$$r_k = \pi^{-1/2} \left(\Gamma(1 + \frac{n}{2}) \text{vol}(\mathcal{D}) \frac{\sigma N^{n-1} \log k}{k} \right)^{1/n} \quad (2)$$



BAMLM with LHS

Critical distance for uniform sampling:

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Critical distance for Latin hypercube sampling:

$$r_k = \pi^{-1/2} \left(\Gamma(1 + \frac{n}{2}) \text{vol}(\mathcal{D}) \frac{\sigma N^{n-1} \log k}{k} \right)^{1/n} \quad (2)$$

Theorem

If r_k is defined by (2) with $\sigma > 4$, even if the sampling continues forever, the total number of local runs started by BAMLM (or AAMLM) is finite almost surely.



Does LHS help?

Problem setup:

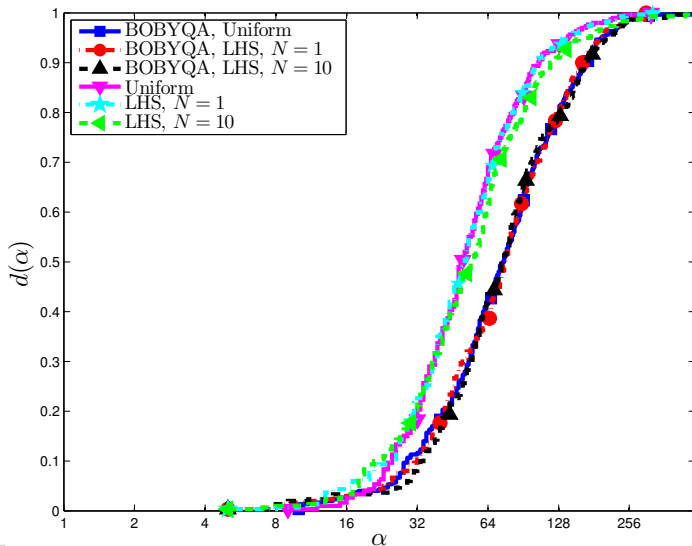
- ▶ 10 different GKLS problems
- ▶ 5 different seeds
- ▶ $n = 2, \dots, 7$

- ▶ Same starting LHS sample of $10n$ points (except for uniform)
- ▶ Same (uniform) r_k value



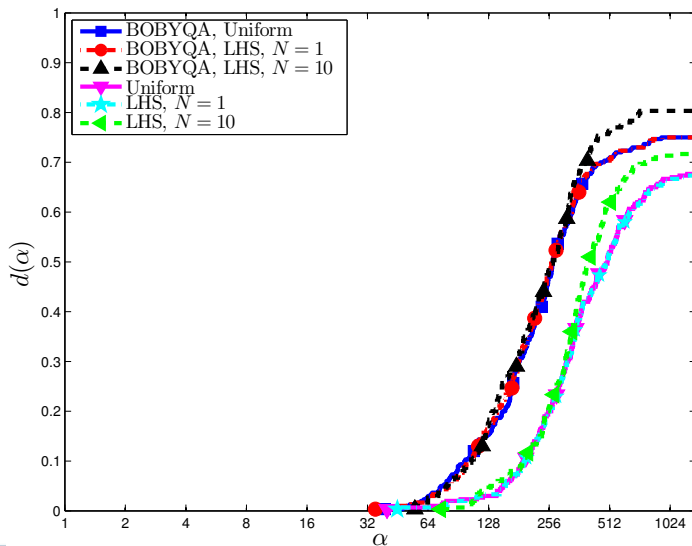
Data Profiles

Within $\sqrt{\frac{10^{-2}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima



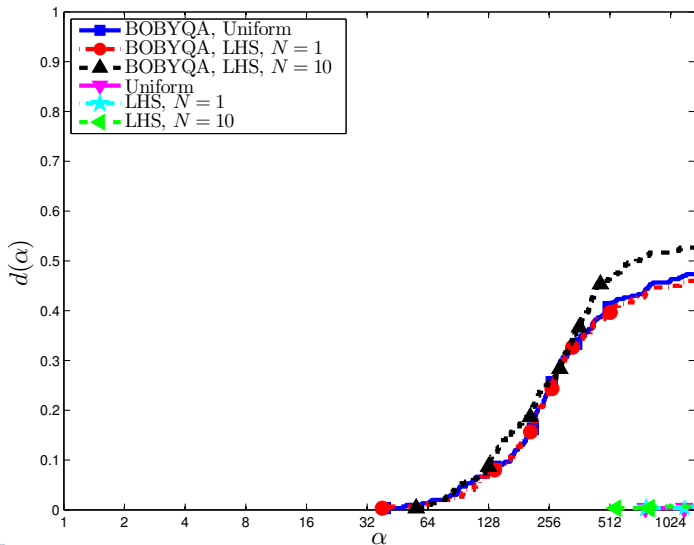
Data Profiles

Within $\sqrt[n]{\frac{10^{-3}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima



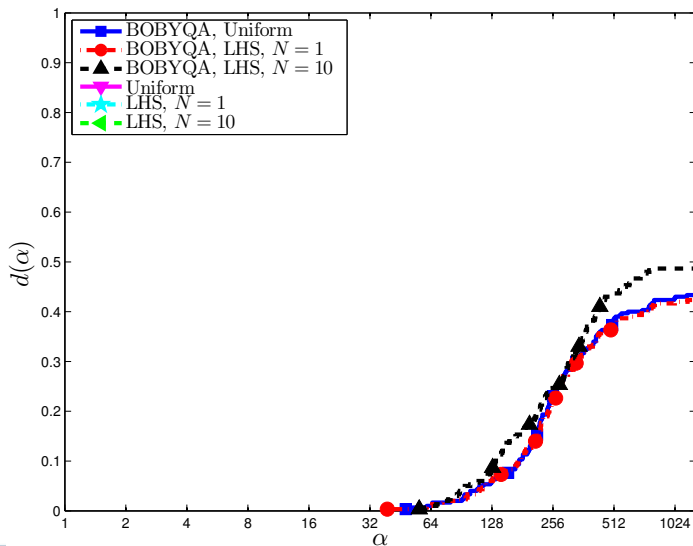
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima



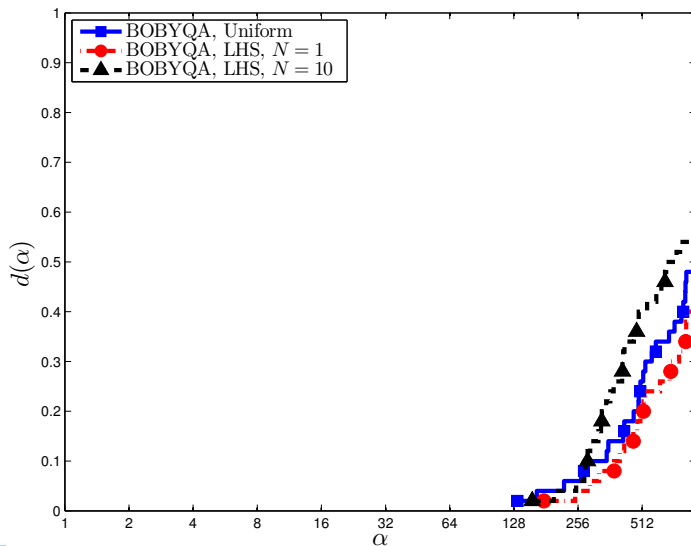
Data Profiles

Within $\sqrt[n]{\frac{10^{-5}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima



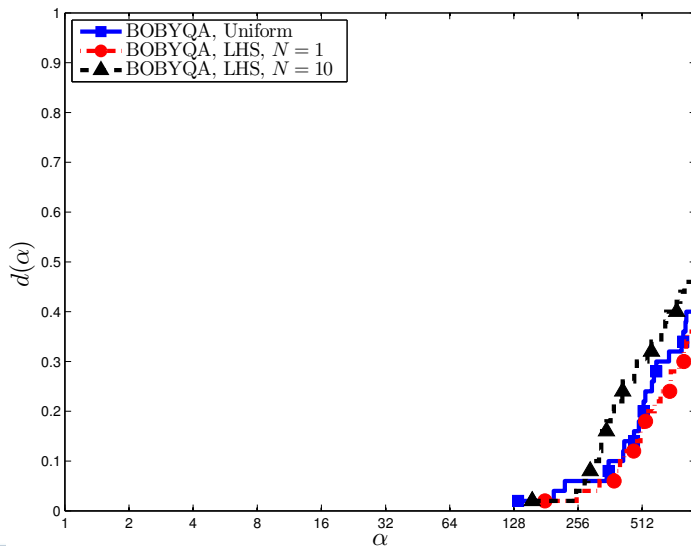
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima, $n = 6$



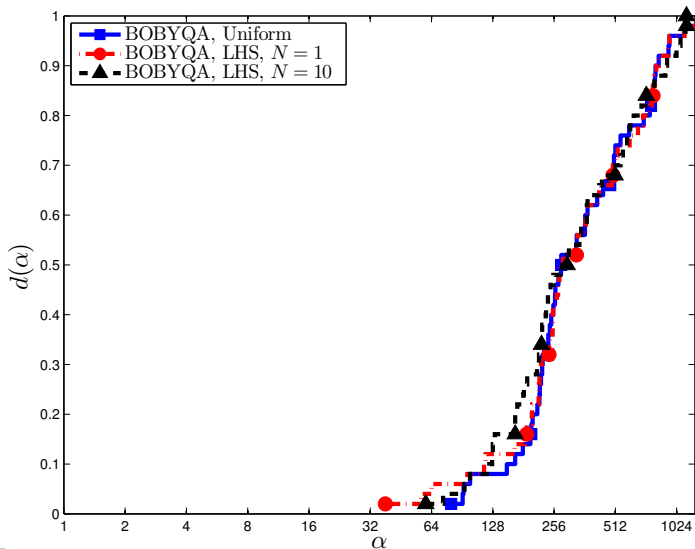
Data Profiles

Within $\sqrt[n]{\frac{10^{-5}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima, $n = 6$



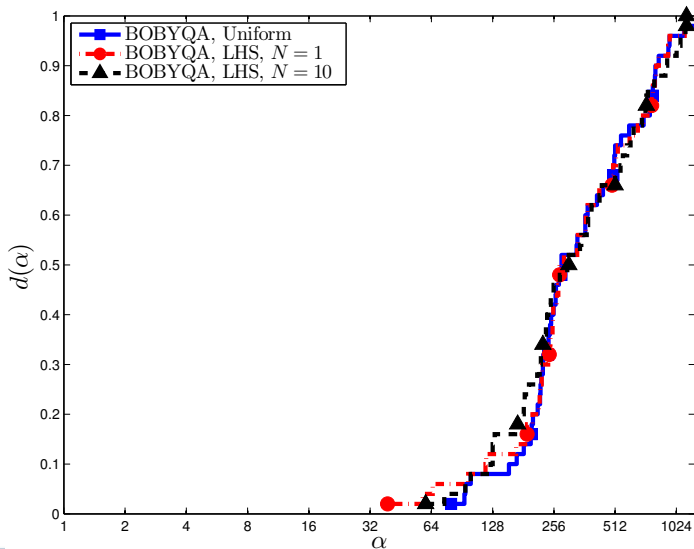
Data Profiles

Within $\sqrt[n]{\frac{10^{-4}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima, $n = 3$



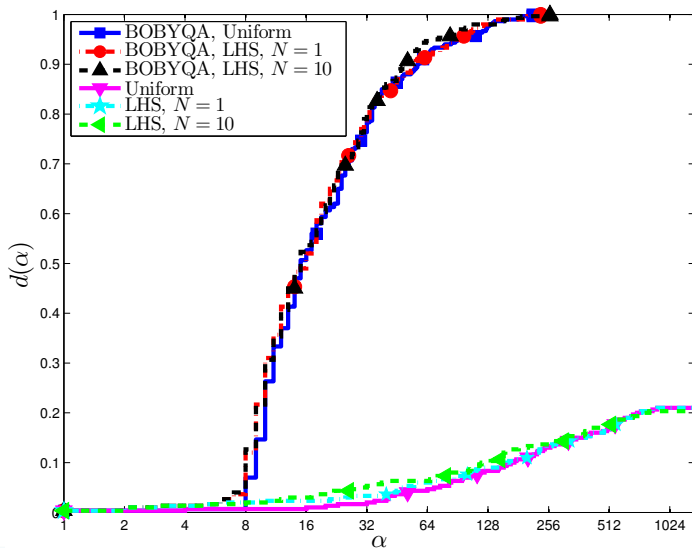
Data Profiles

Within $\sqrt[n]{\frac{10^{-5}\Gamma(\frac{n}{2}+1)}{\pi^{n/2}}}$ of 5 best minima, $n = 3$



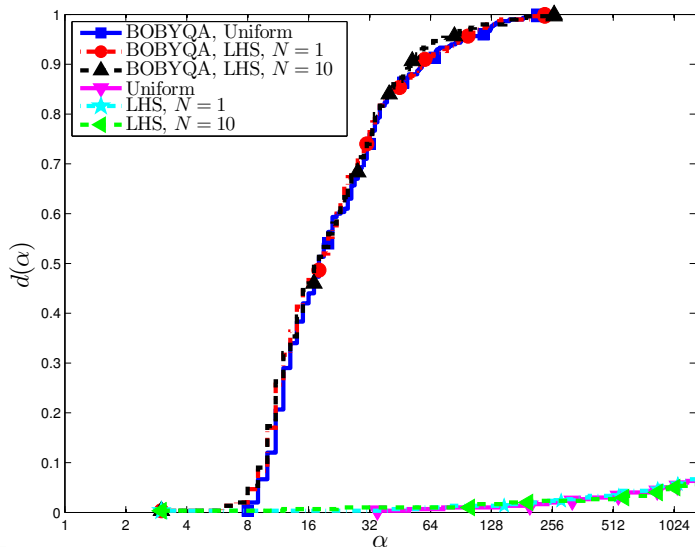
Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-2}) \left(f(x_0) - f_{(1)}^* \right)$$



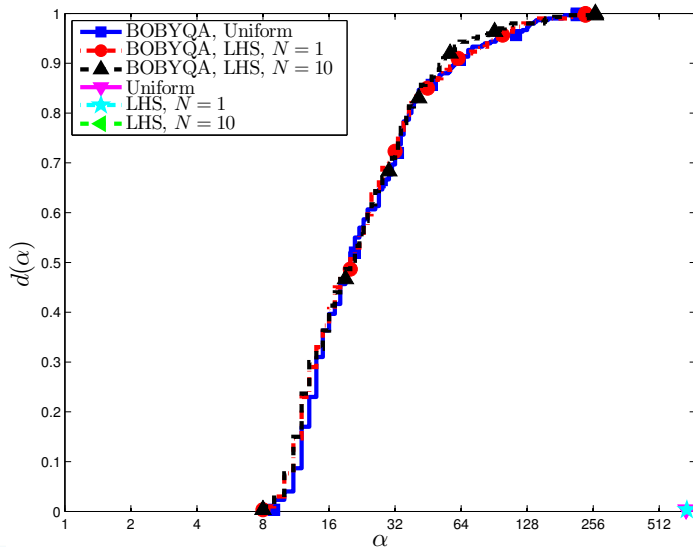
Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-3}) \left(f(x_0) - f_{(1)}^* \right)$$



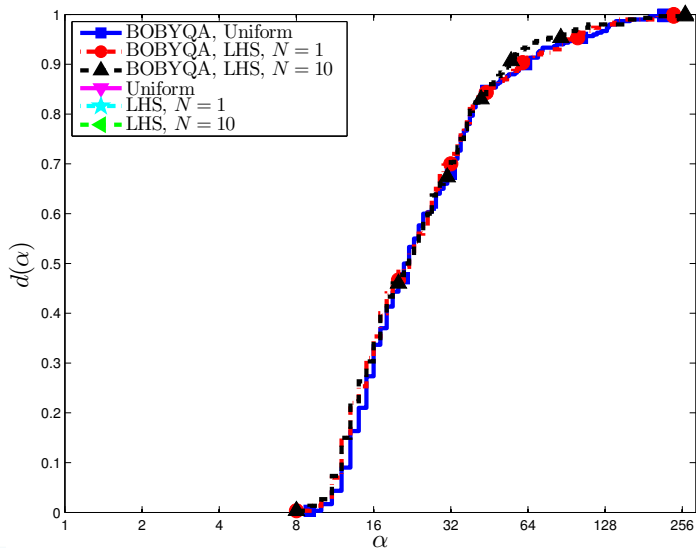
Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-4}) \left(f(x_0) - f_{(1)}^* \right)$$



Data Profiles

$$f(x) - f_{(1)}^* \leq (1 - 10^{-5}) \left(f(x_0) - f_{(1)}^* \right)$$



Closing Remarks

- ▶ Concurrent function evaluations can locate multiple minima while efficiently finding a global minimum.



Closing Remarks

- ▶ Concurrent function evaluations can locate multiple minima while efficiently finding a global minimum.
- ▶ Latin hypercube sampling appears to help find more minima in higher-dimensional problems.

Questions:

- ▶ Finding (or designing) the best local solver for our framework?
- ▶ Best way to process the queue?



Algorithm 3: AAML

Give each worker a point to evaluate

for $k = 1, 2, \dots$ **do**

 Receive from (longest waiting) worker w that has evaluated f

 Update \mathcal{H}_k and r_k

if *point evaluated by w is from an active run* **then**

if *Run is complete* **then**

 Update X_k^* , and mark points inactive

else

 Add the next point in its localopt run (not in \mathcal{H}_k) to Q_L

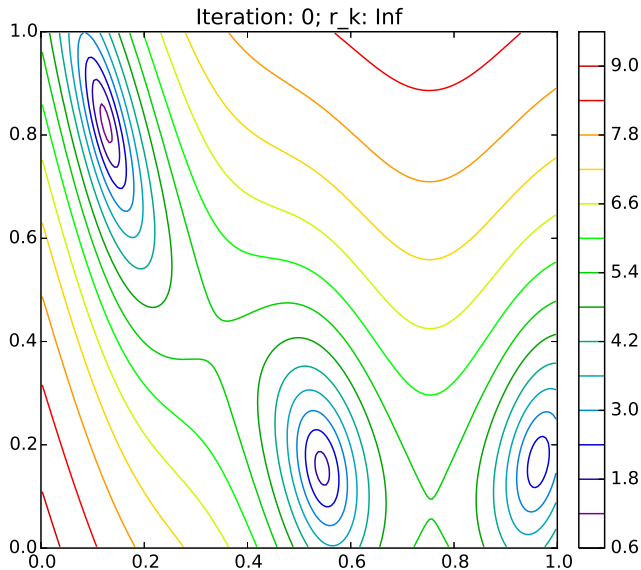
 Start run(s) at all point(s) satisfying (S1)–(S4), (L1)–(L6)

 Add the subsequent point (not in \mathcal{H}_k) from each run to Q_L

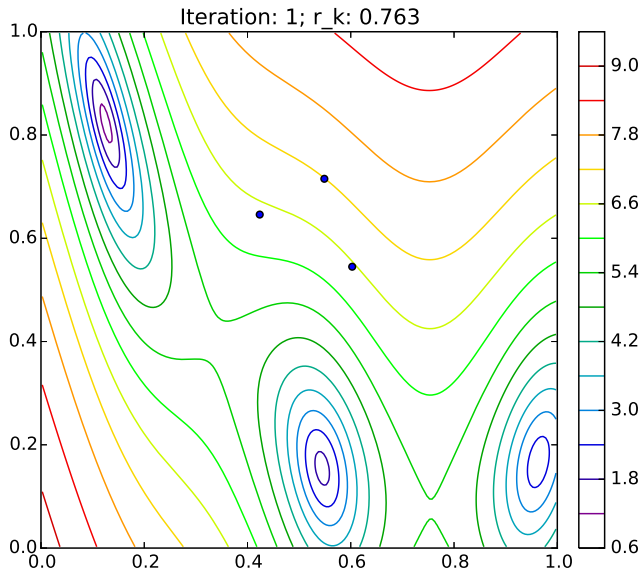
 Merge runs in Q_L with candidate minima within 2ν of each other

 Give w a point at which to evaluate f , either from Q_L or \mathcal{R}

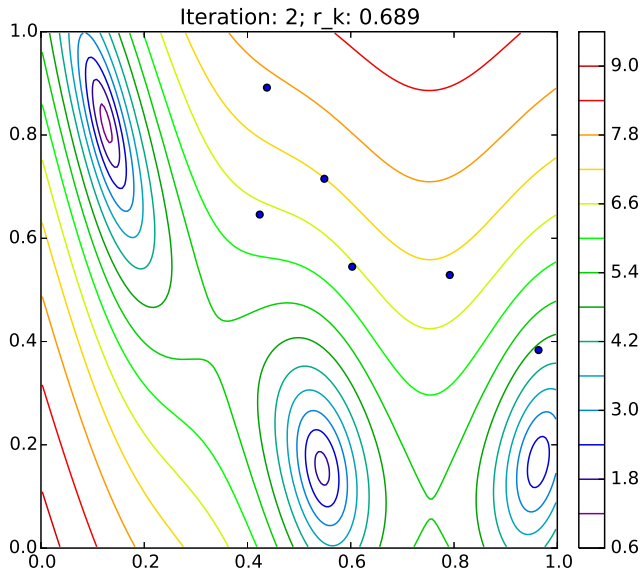
Pausing runs



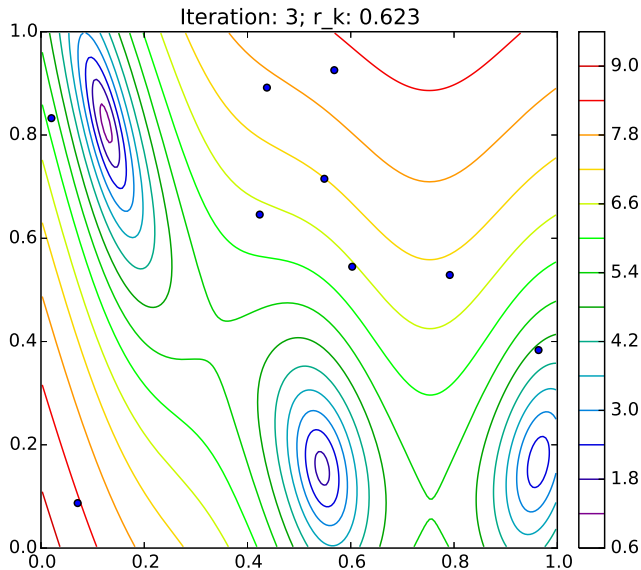
Pausing runs



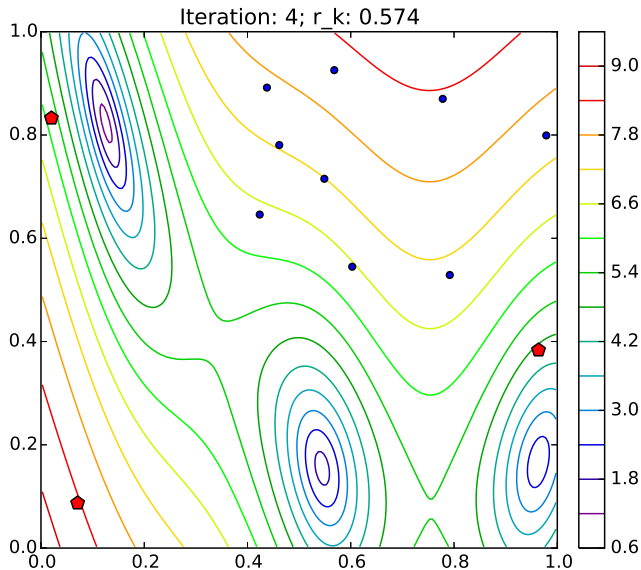
Pausing runs



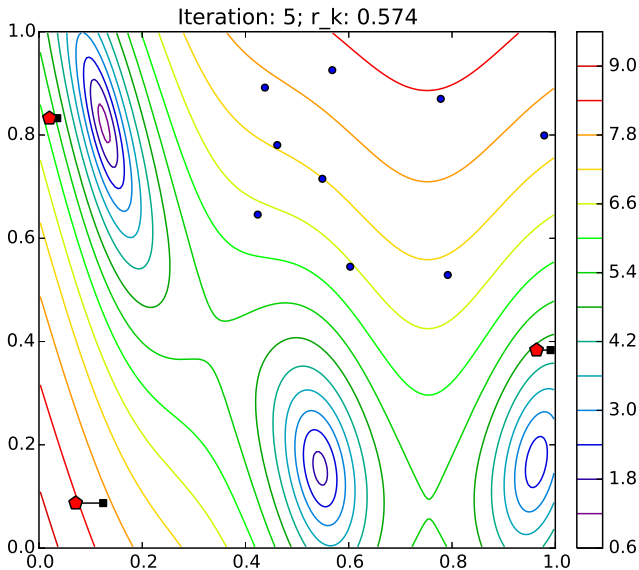
Pausing runs



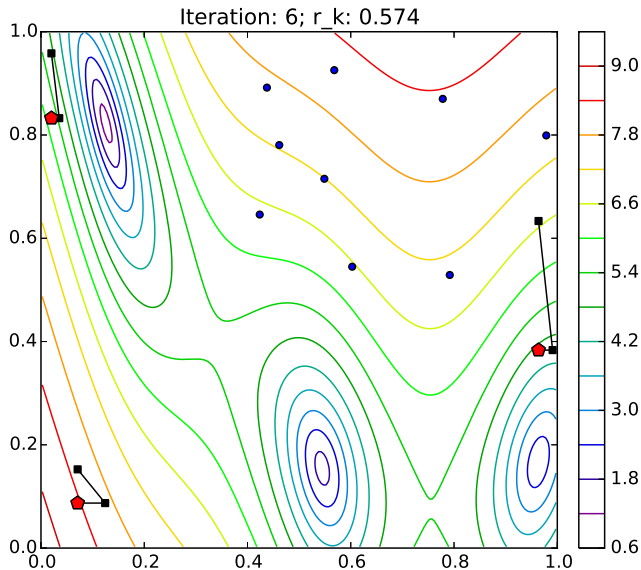
Pausing runs



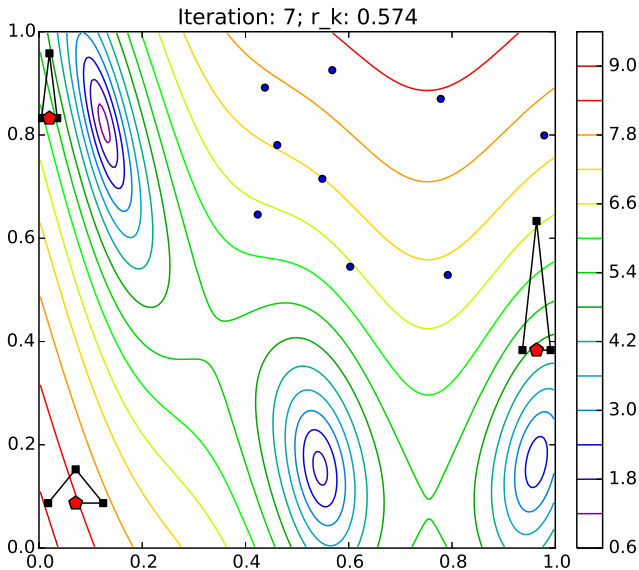
Pausing runs



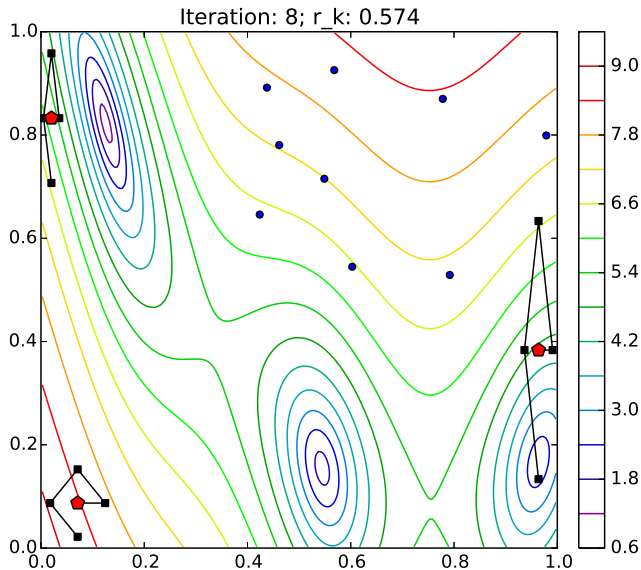
Pausing runs



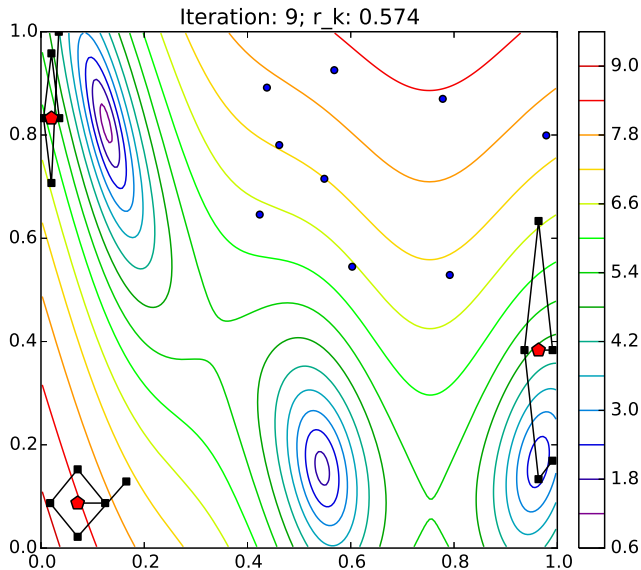
Pausing runs



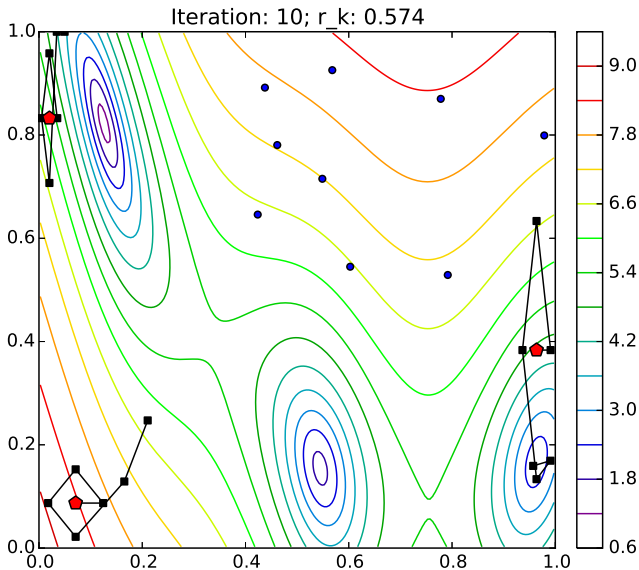
Pausing runs



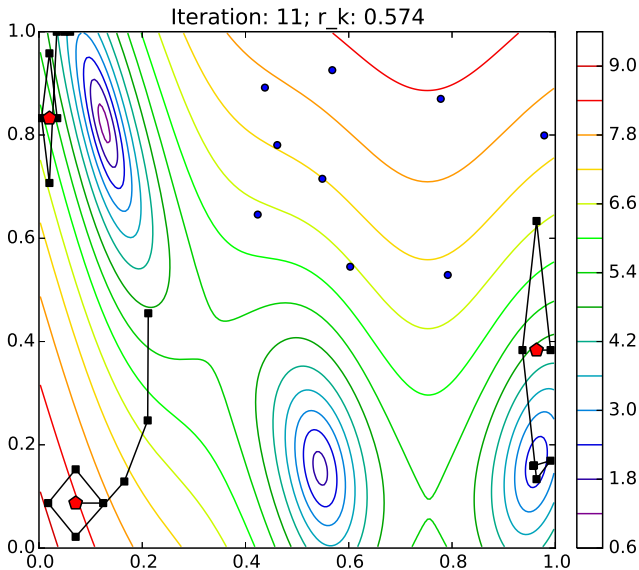
Pausing runs



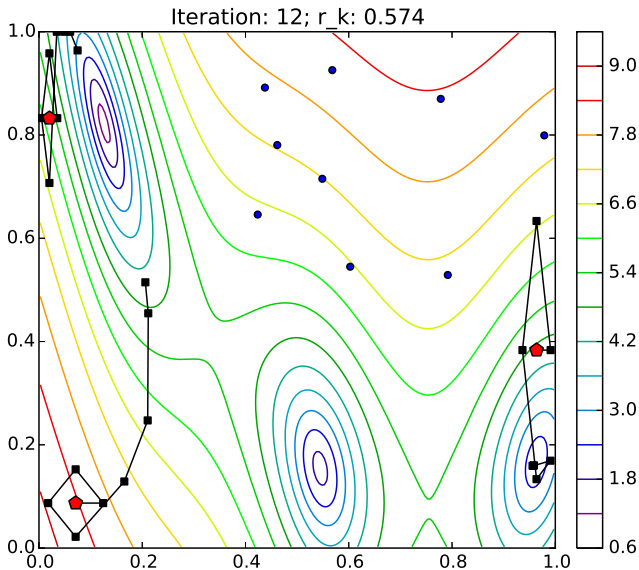
Pausing runs



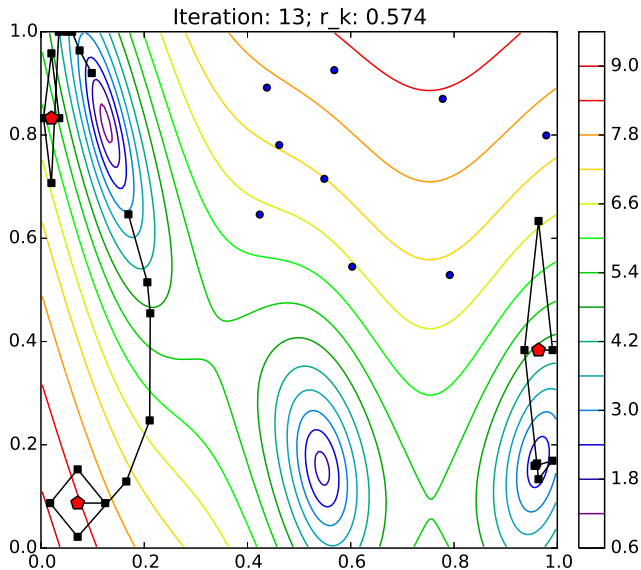
Pausing runs



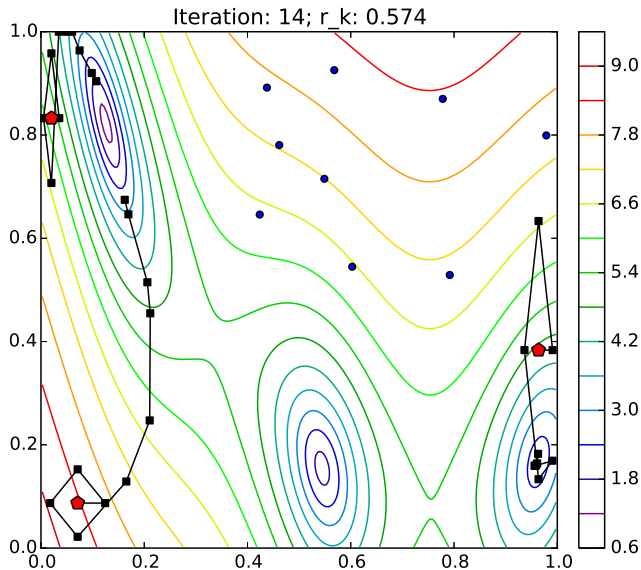
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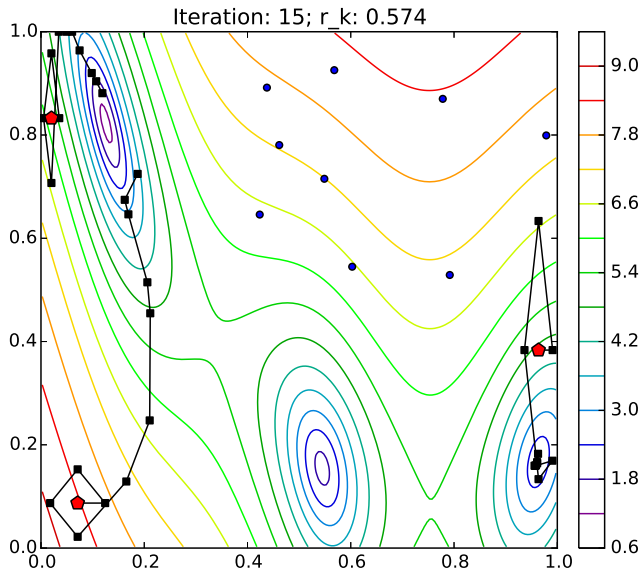
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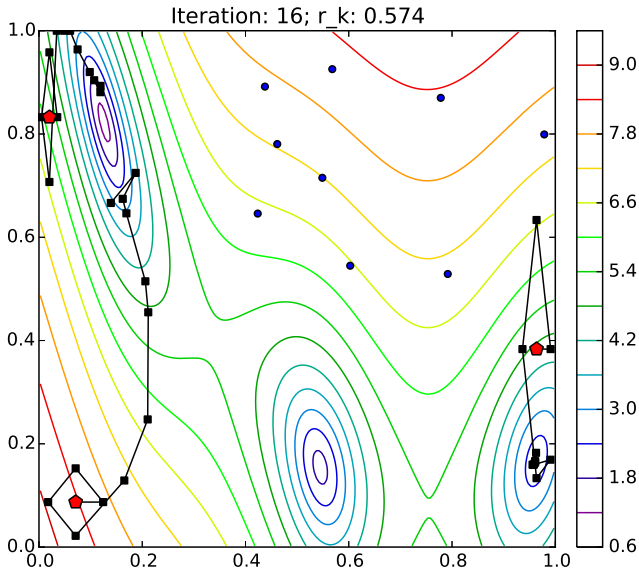
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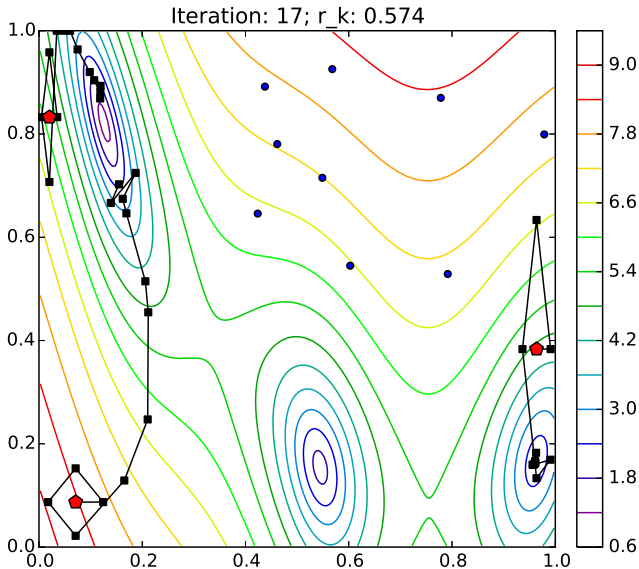
Pausing runs



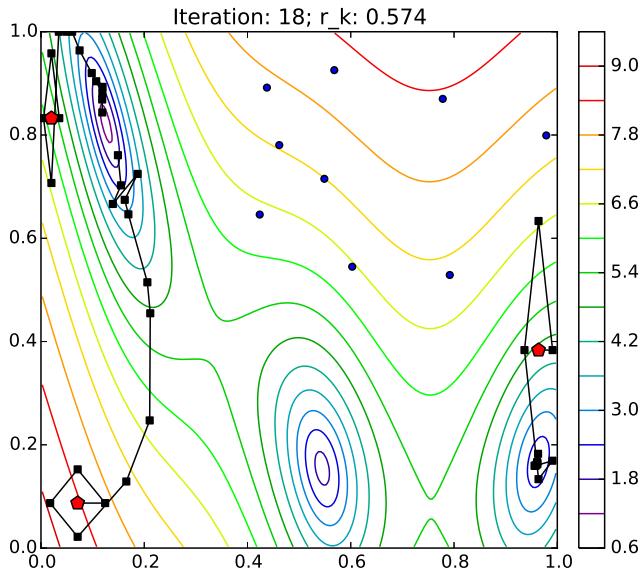
Pausing runs



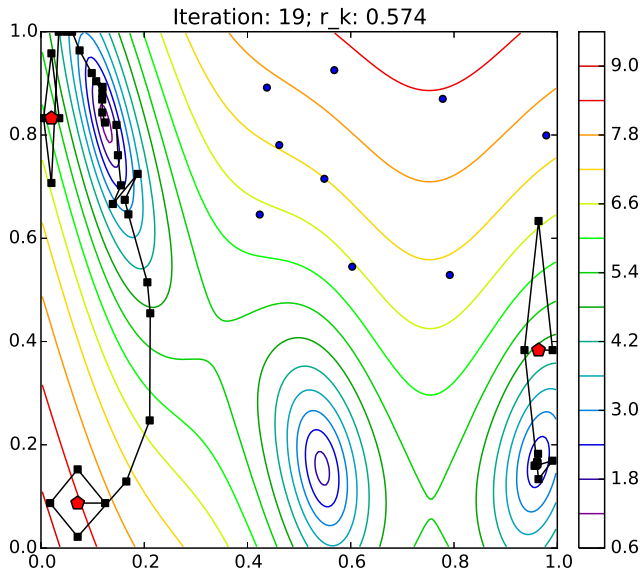
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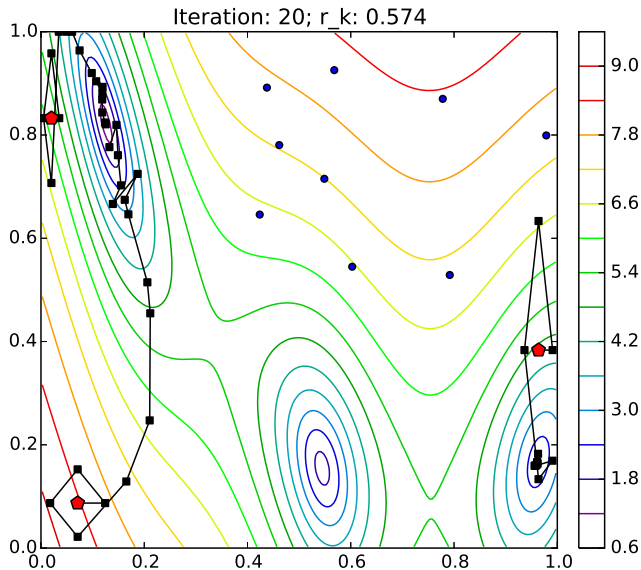
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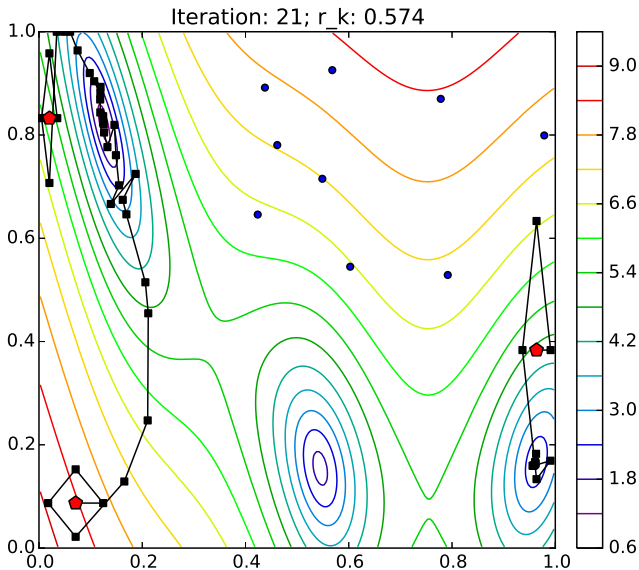
Pausing runs



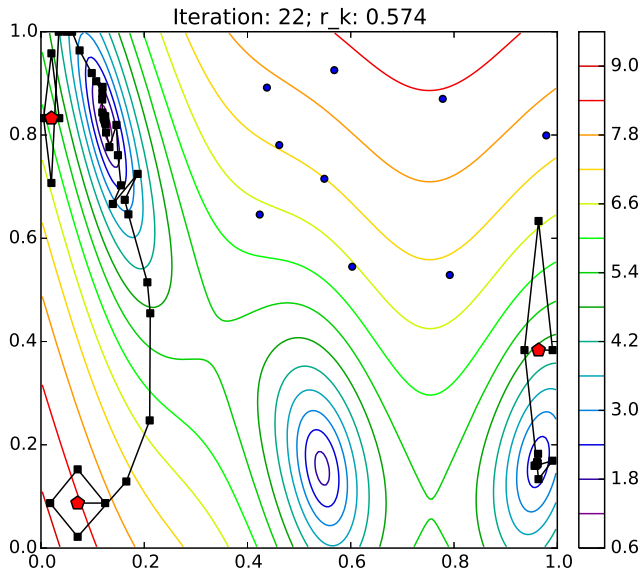
Pausing runs



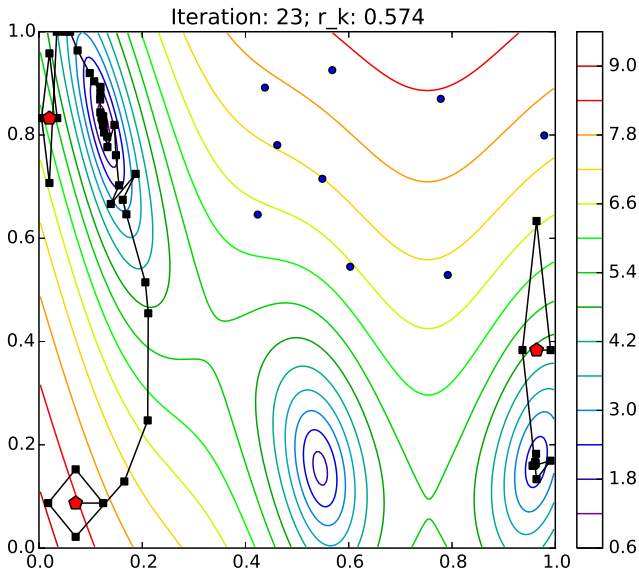
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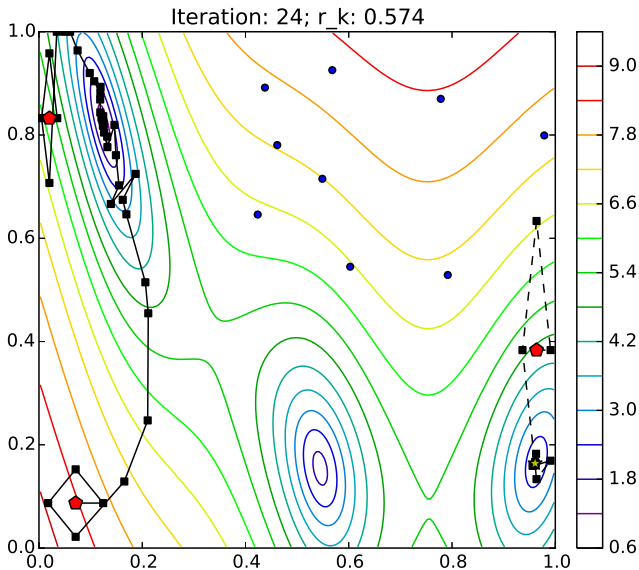
Pausing runs



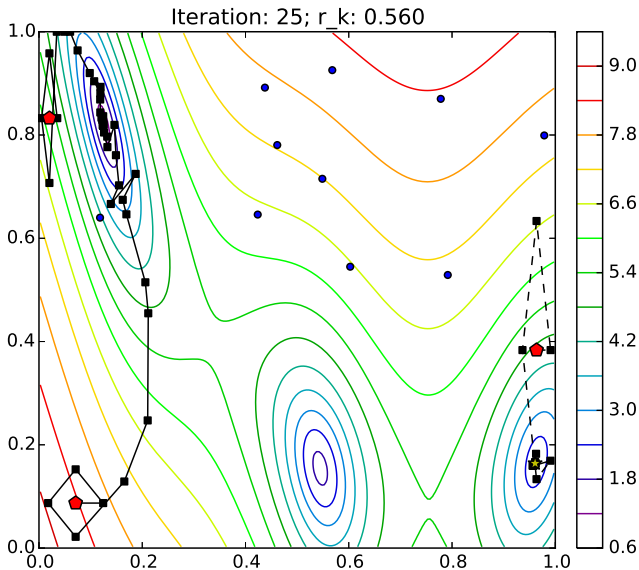
Pausing runs



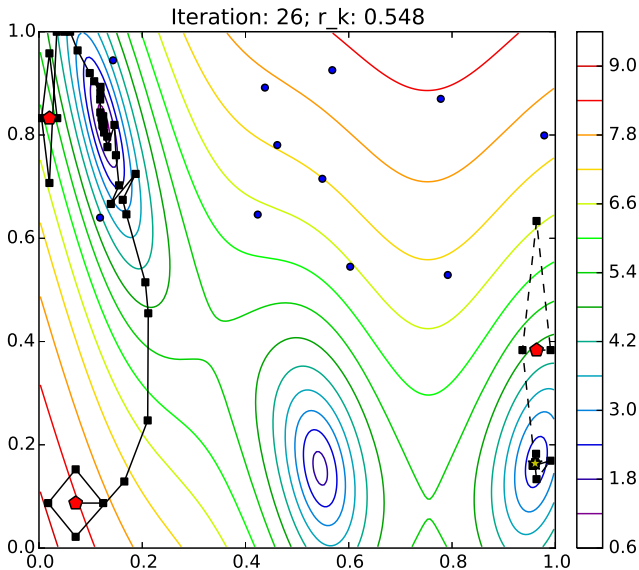
Pausing runs



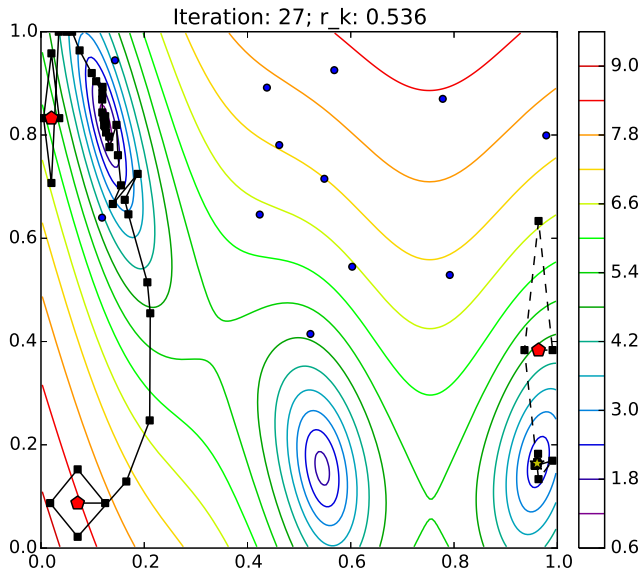
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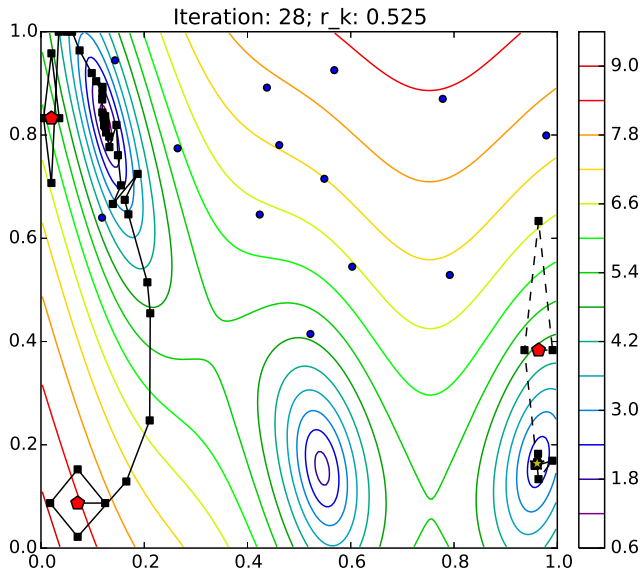
Pausing runs



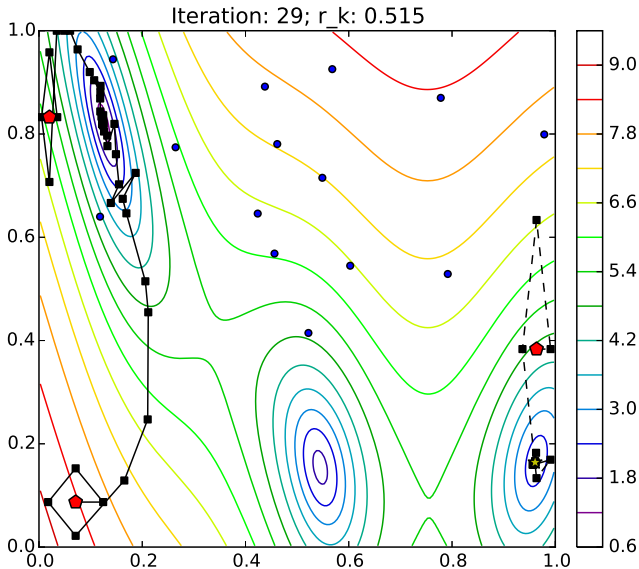
Pausing runs



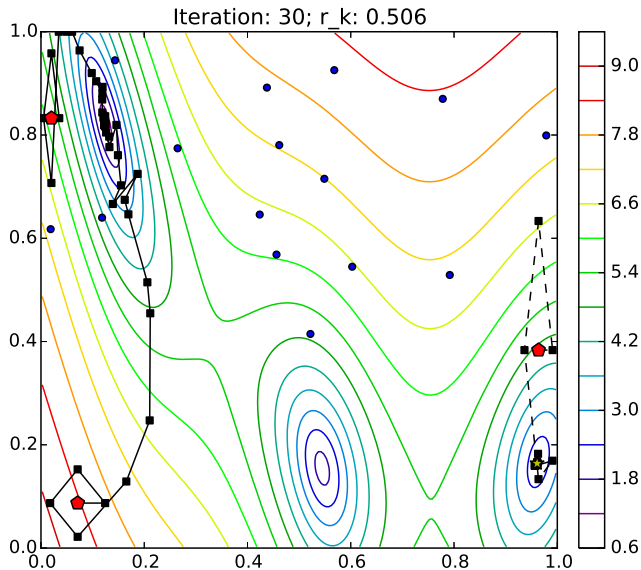
Pausing runs



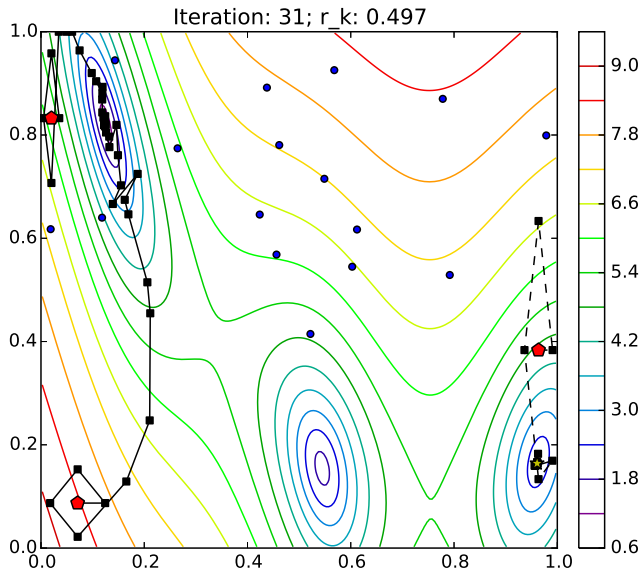
Pausing runs



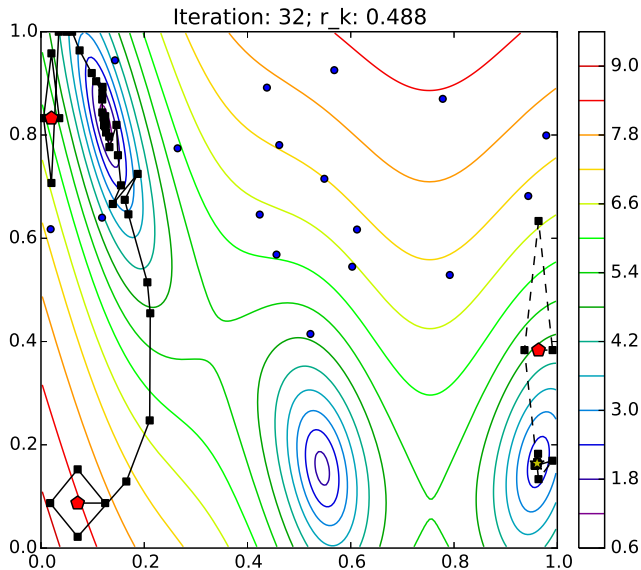
Pausing runs



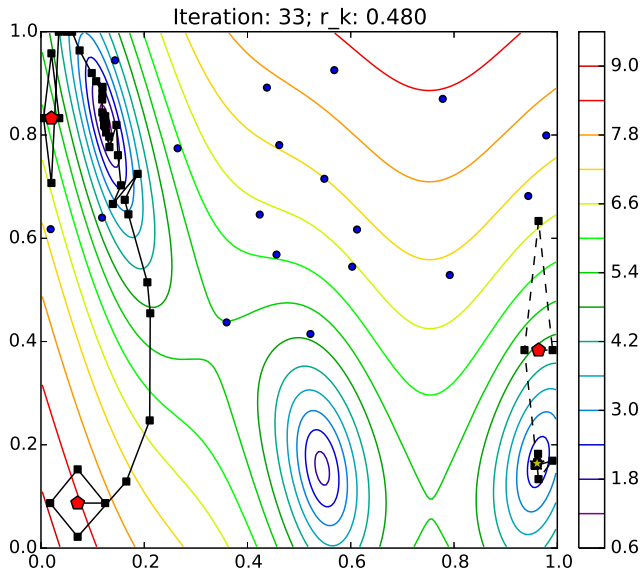
Pausing runs



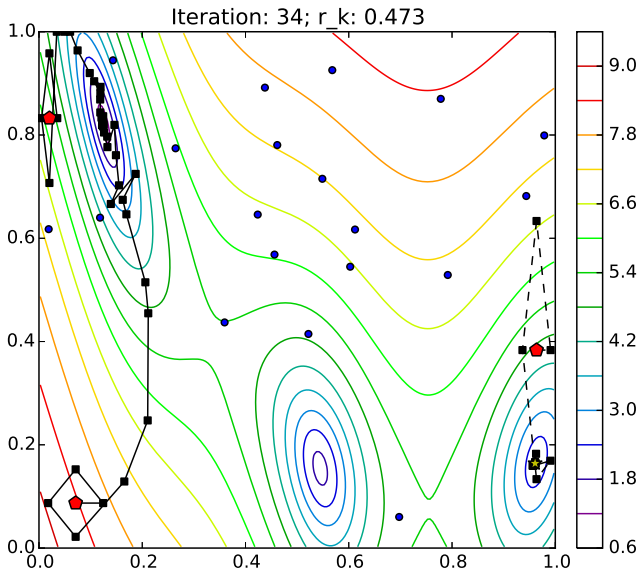
Pausing runs



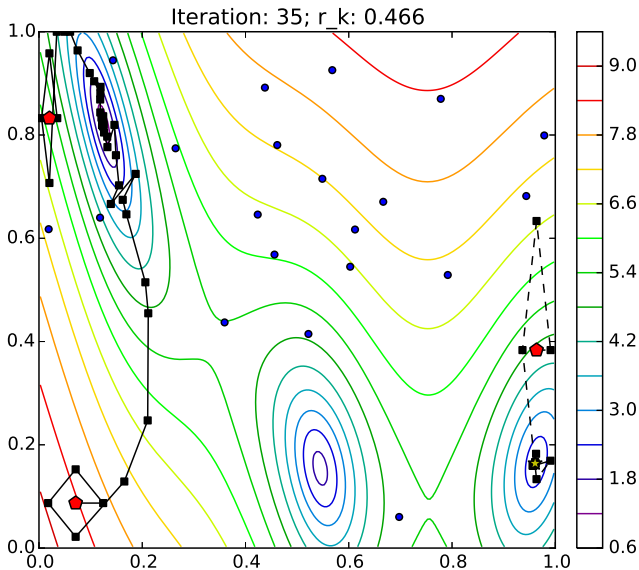
Pausing runs



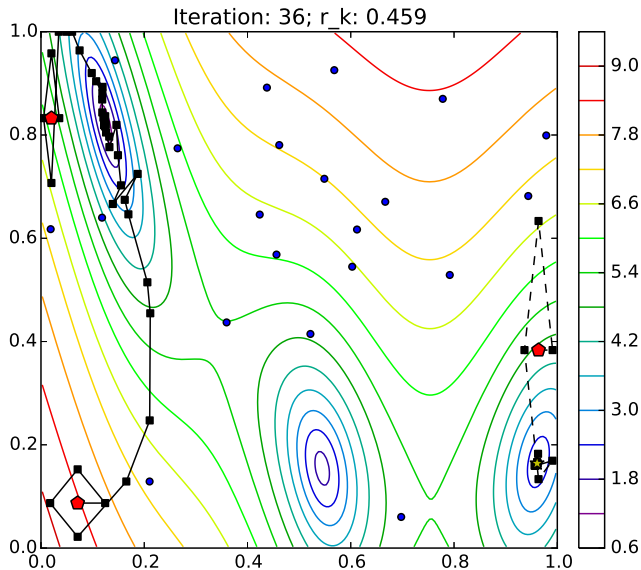
Pausing runs



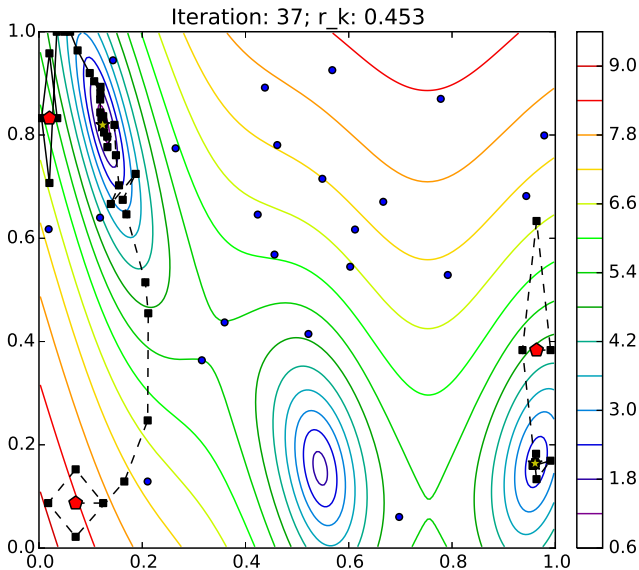
Pausing runs



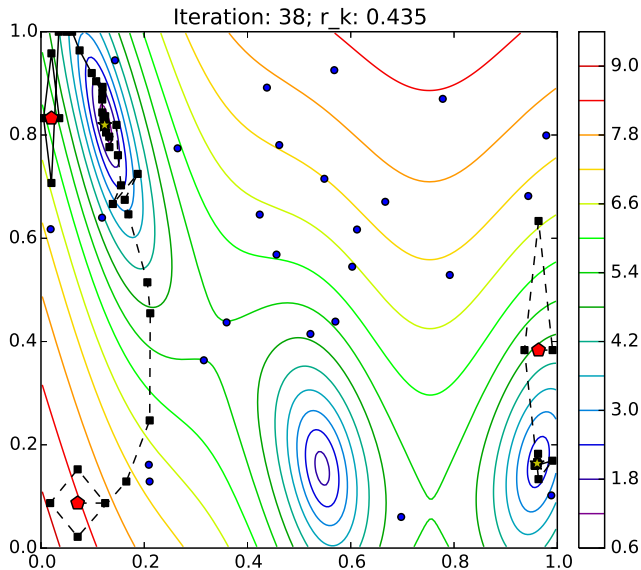
Pausing runs



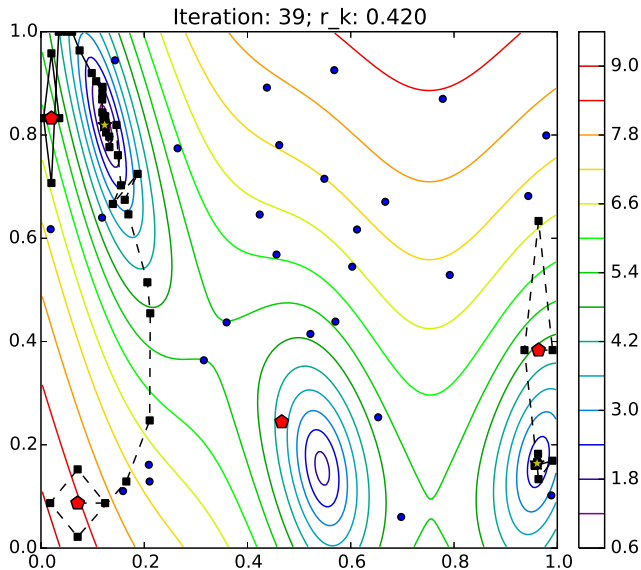
Pausing runs



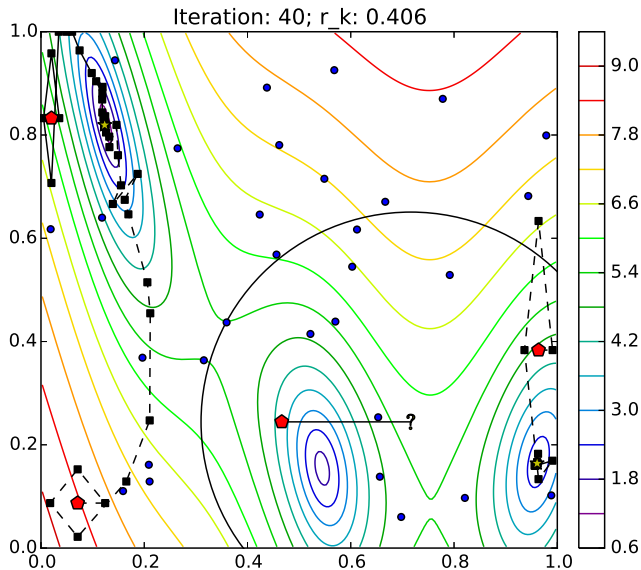
Pausing runs



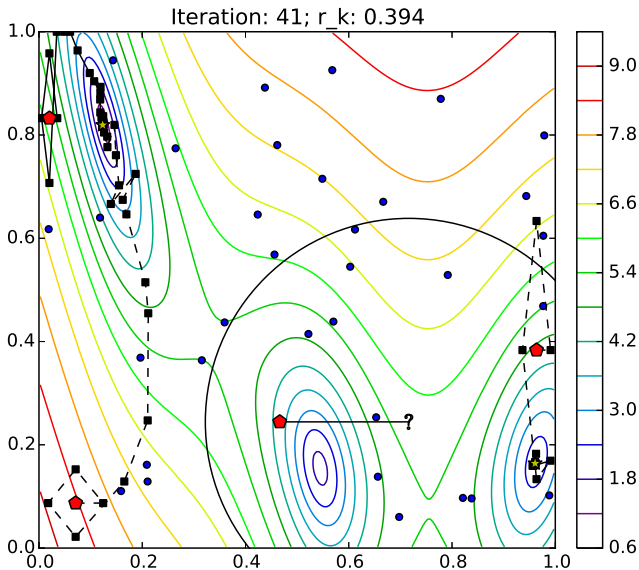
Pausing runs



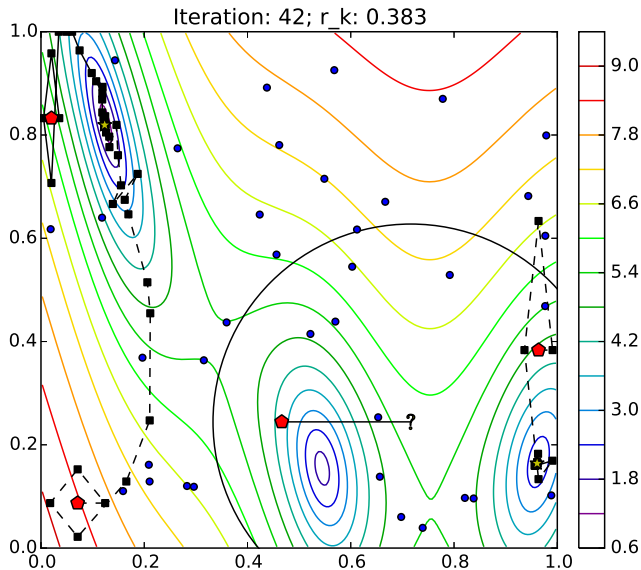
Pausing runs



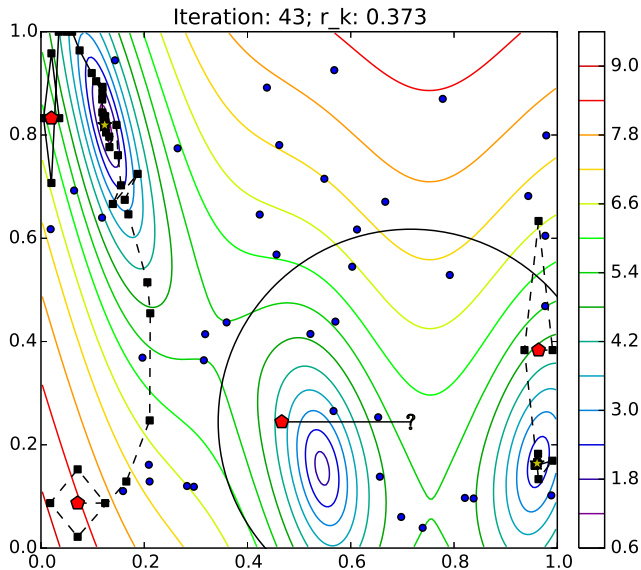
Pausing runs



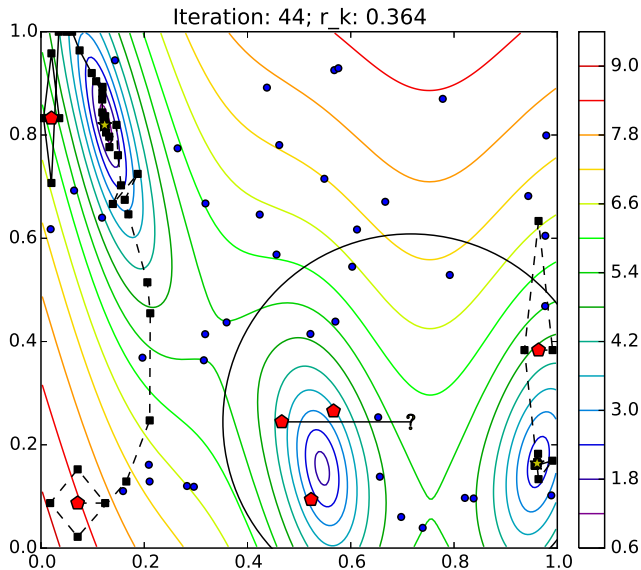
Pausing runs



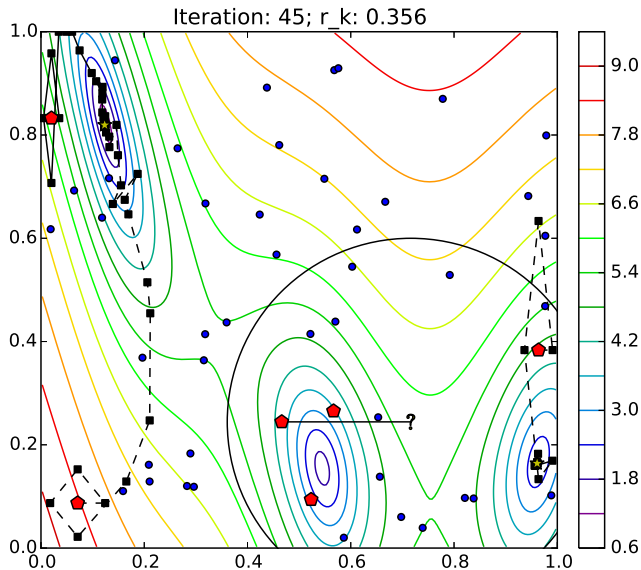
Pausing runs



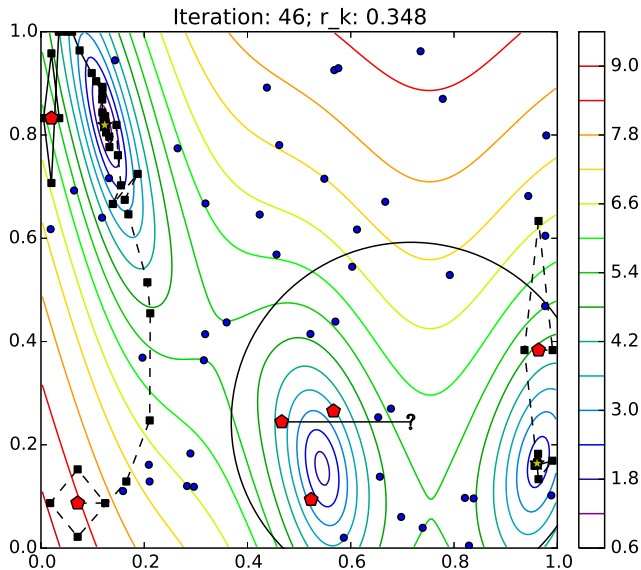
Pausing runs



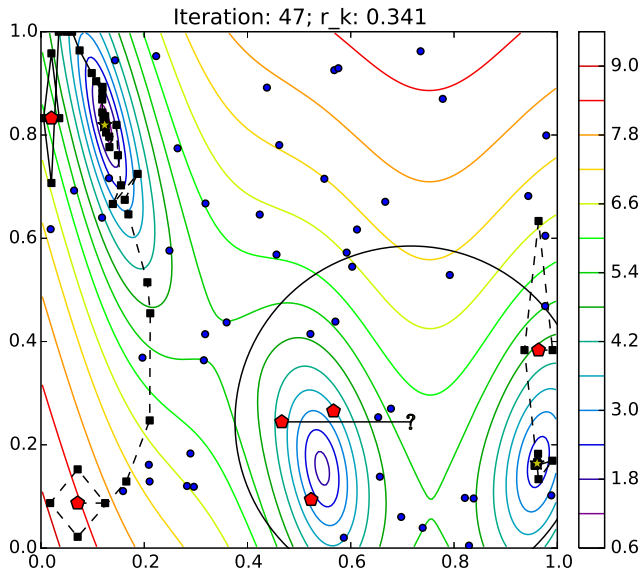
Pausing runs



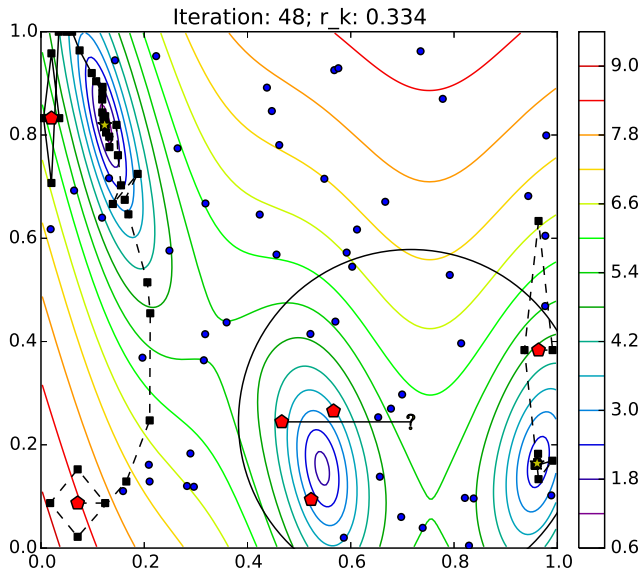
Pausing runs



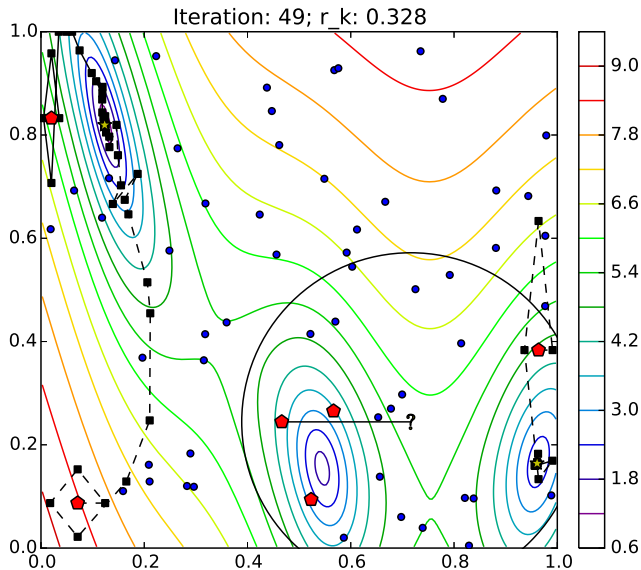
Pausing runs



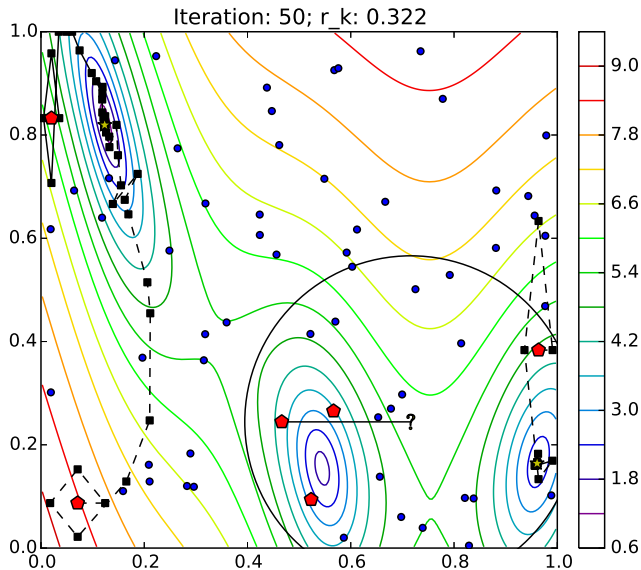
Pausing runs



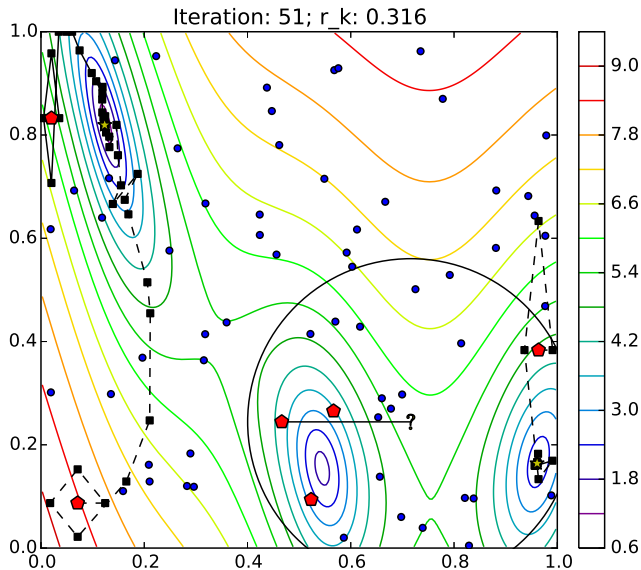
Pausing runs



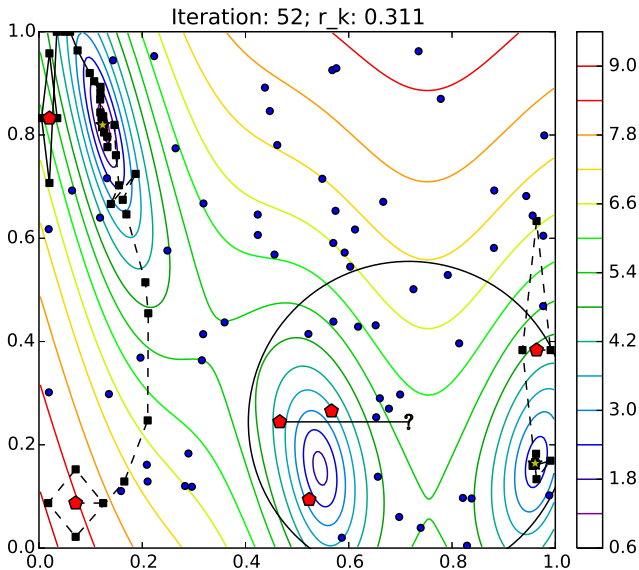
Pausing runs



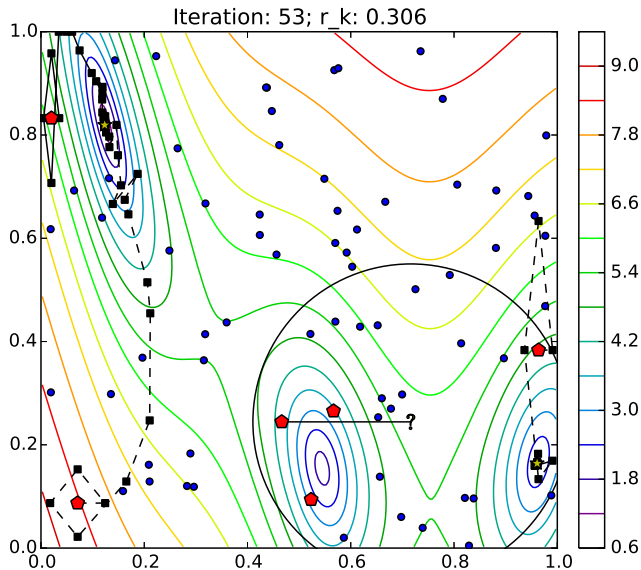
Pausing runs



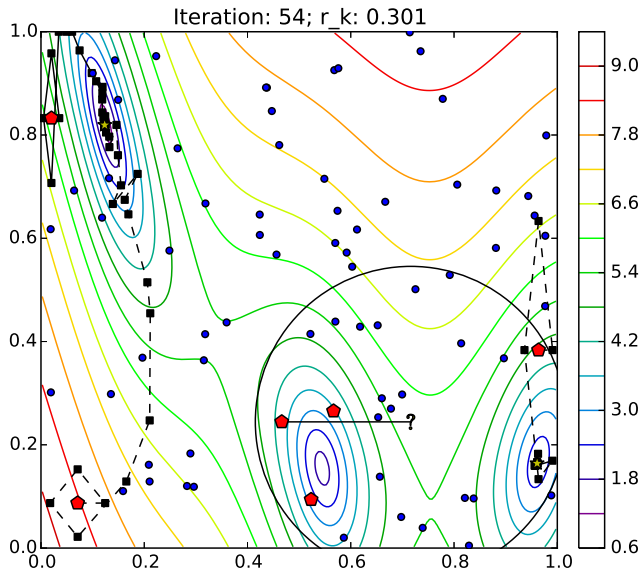
Pausing runs



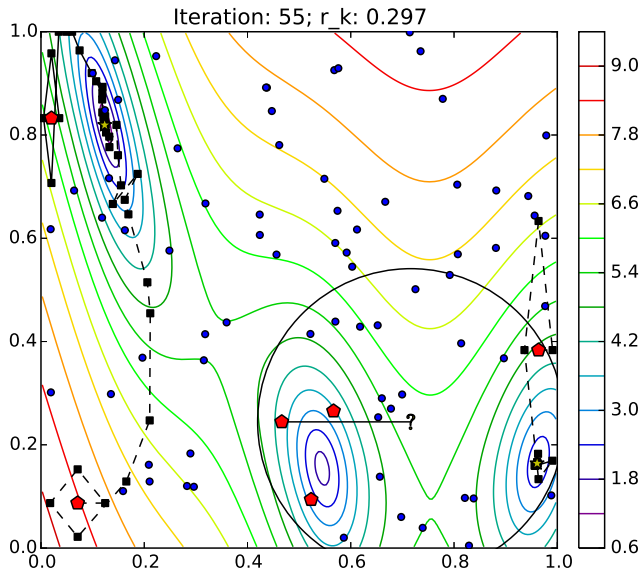
Pausing runs



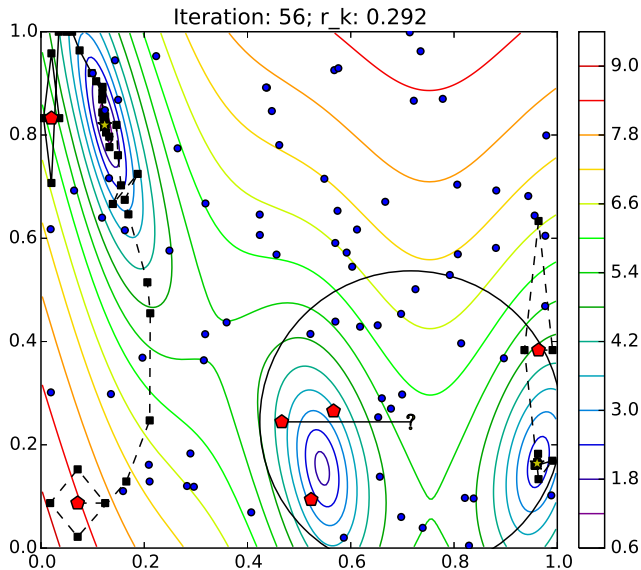
Pausing runs



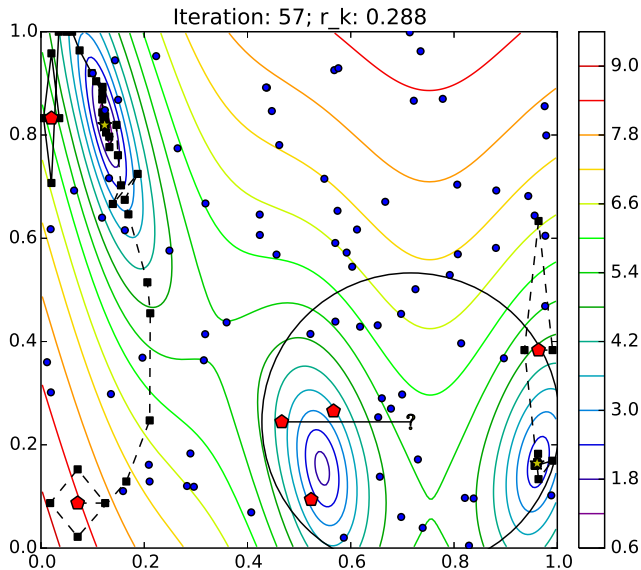
Pausing runs



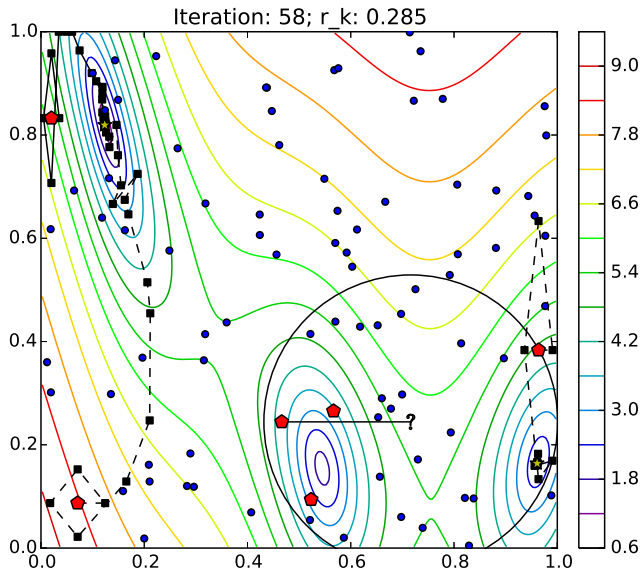
Pausing runs



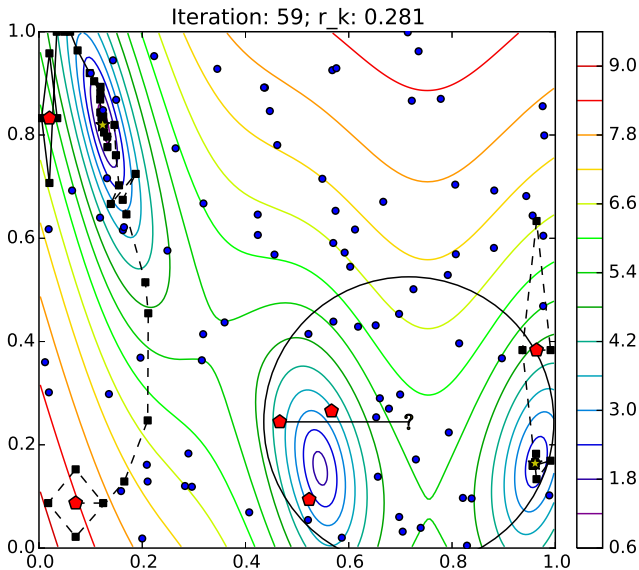
Pausing runs



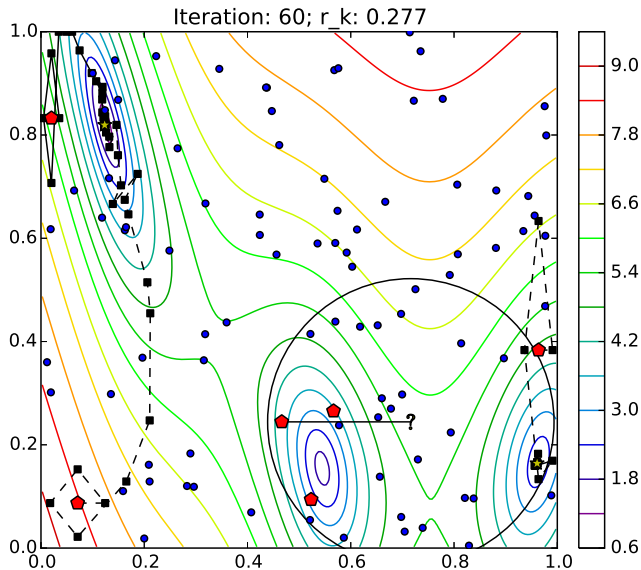
Pausing runs



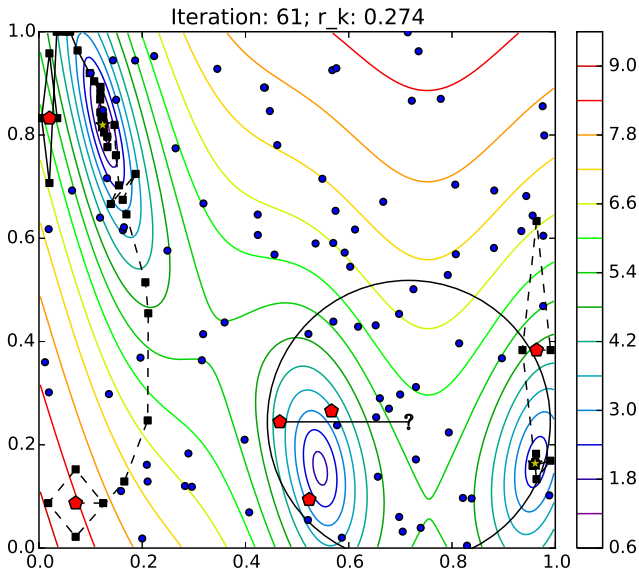
Pausing runs



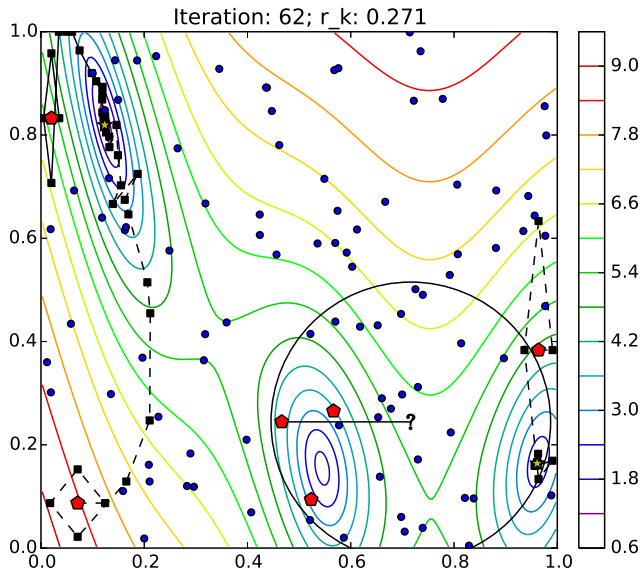
Pausing runs



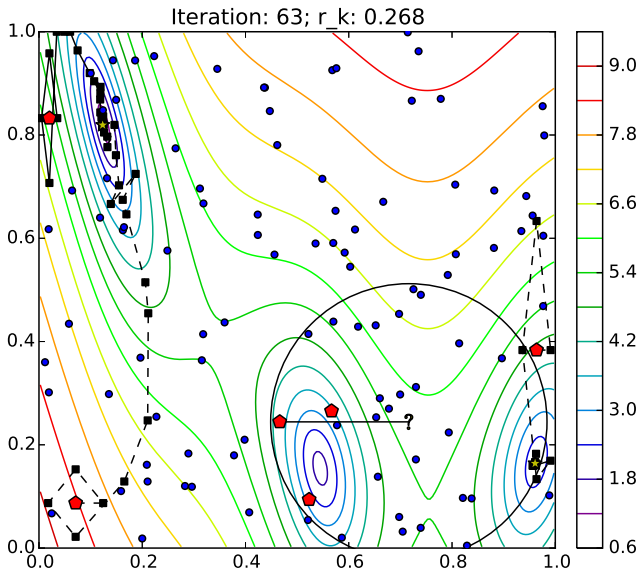
Pausing runs



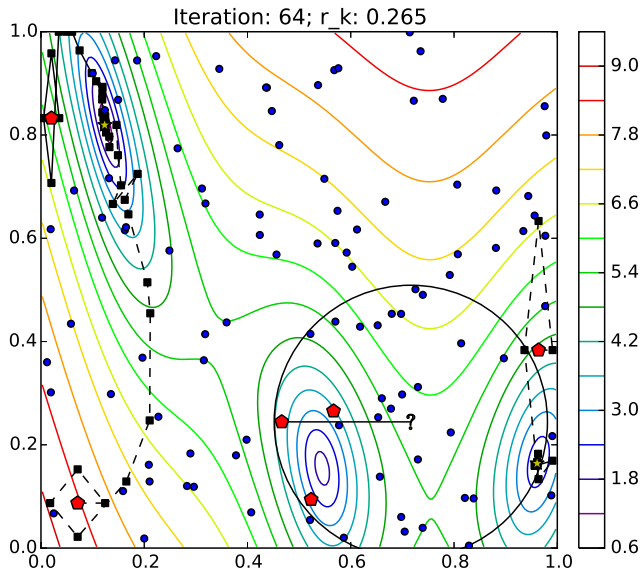
Pausing runs



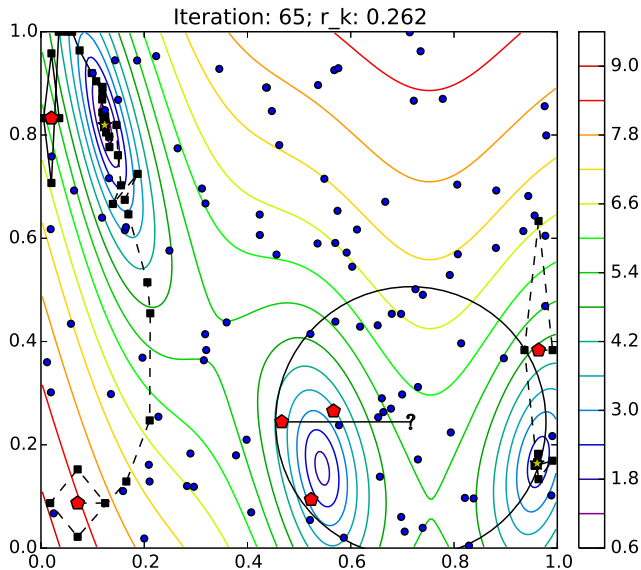
Pausing runs



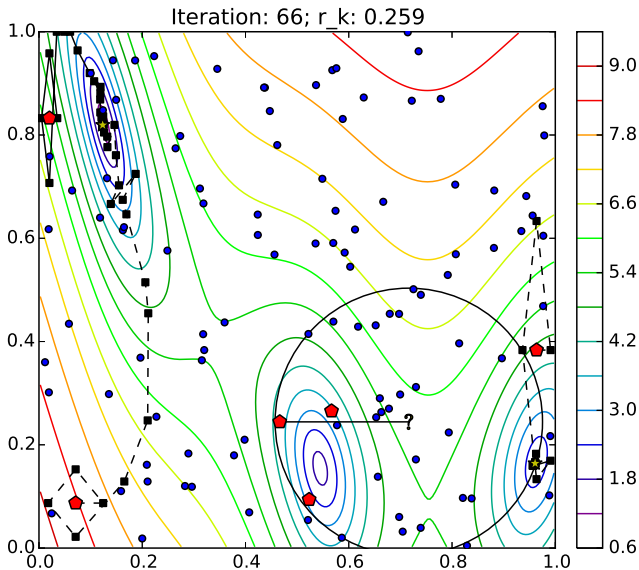
Pausing runs



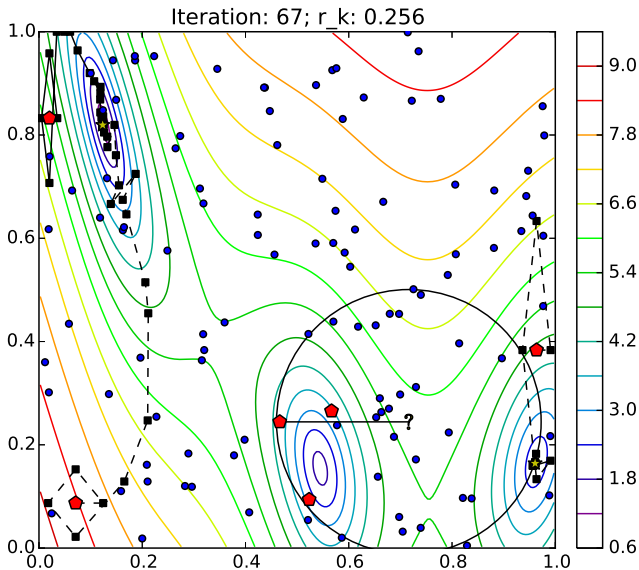
Pausing runs



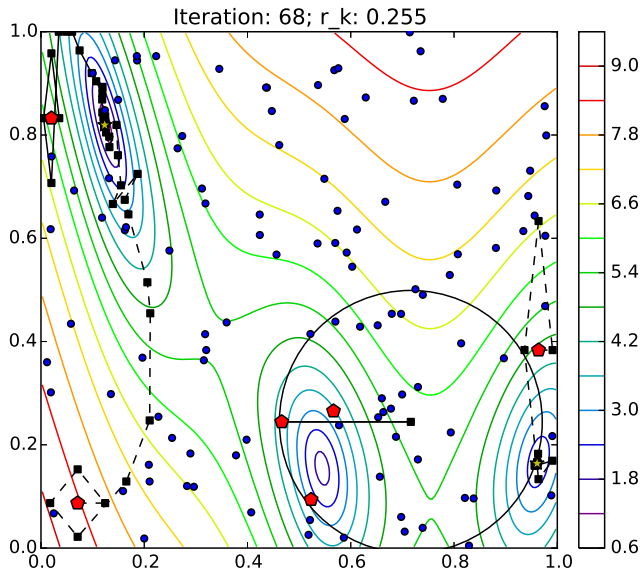
Pausing runs



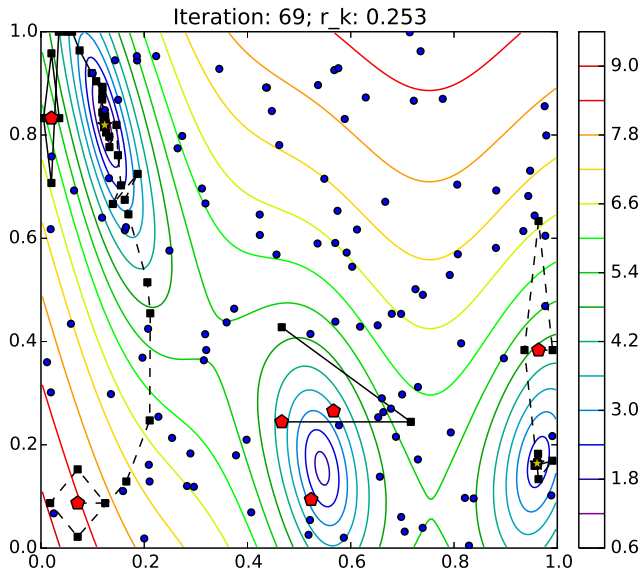
Pausing runs



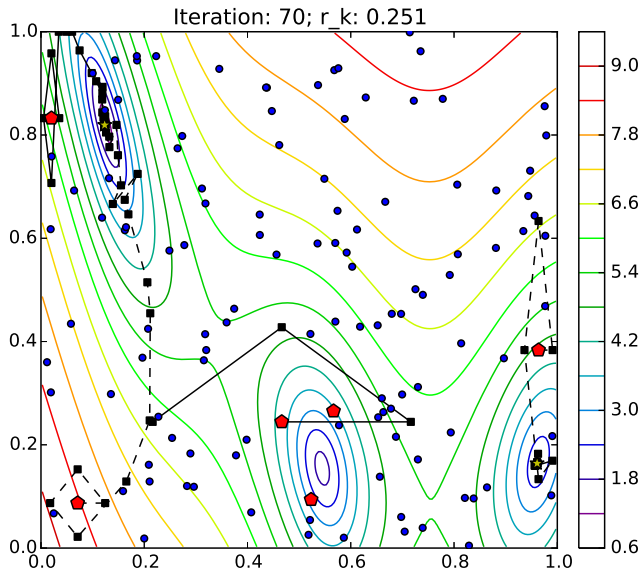
Pausing runs



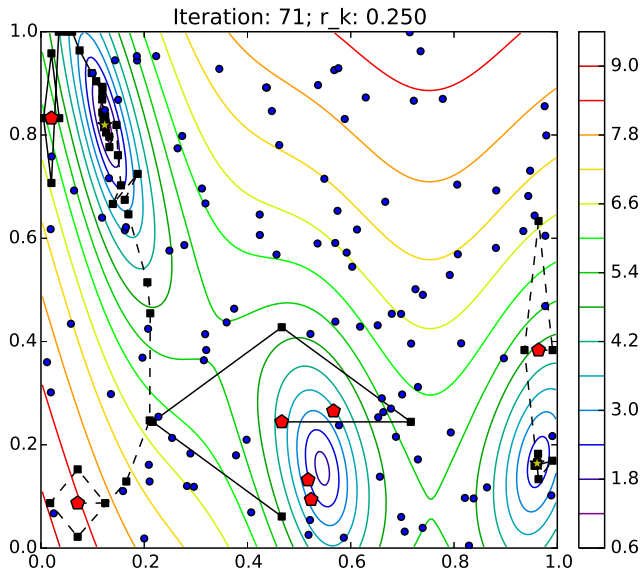
Pausing runs



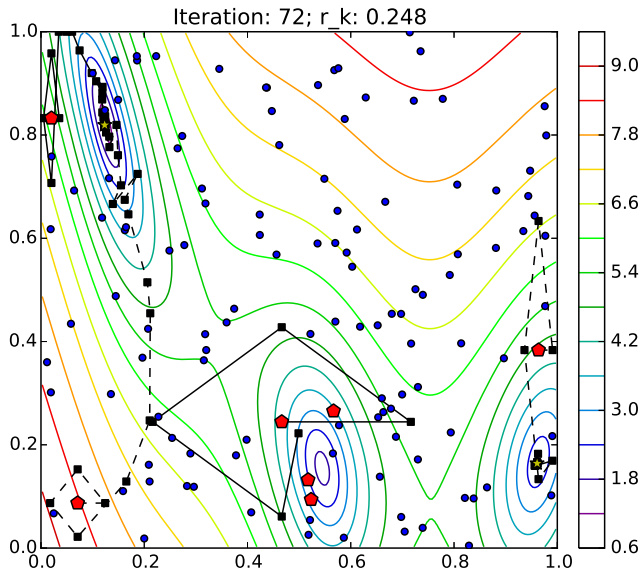
Pausing runs



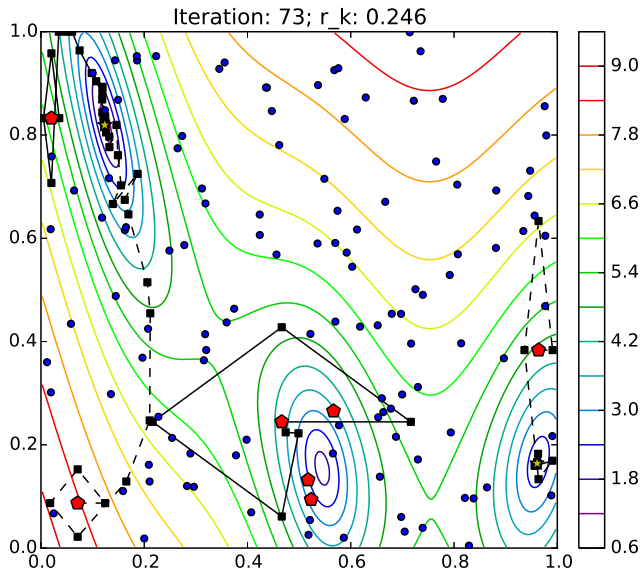
Pausing runs



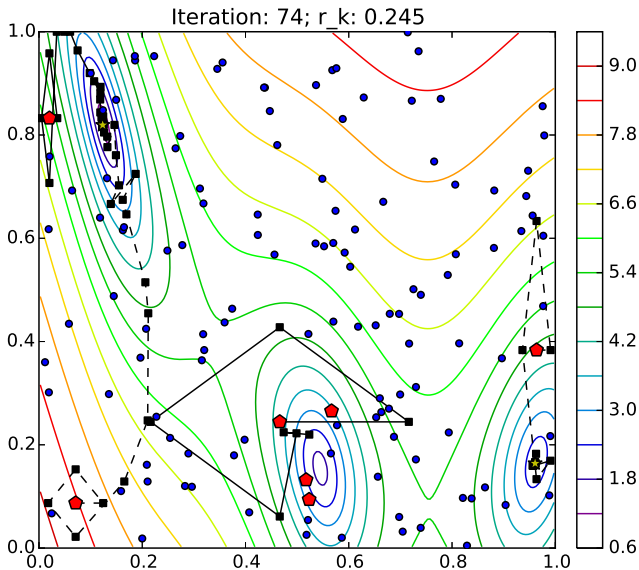
Pausing runs



Pausing runs



Pausing runs



Pausing runs

